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Weighted aggregation of partial rankings using Ant Colony Optimization

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Abstract

The aggregation of preferences (expressed in the form of rankings) from multiple experts is a well-studied topic in a number of fields. The Kemeny ranking problem aims at computing an aggregated ranking having minimal distance to the global consensus. However, it assumes that these rankings will be complete, i.e., all elements are explicitly ranked by the expert. This assumption may not simply hold when, for instance, an expert ranks only the top- K items of interest, thus creating a partial ranking. In this paper we formalize the *weighted Kemeny ranking problem for partial rankings*, an extension of the Kemeny ranking problem that is able to aggregate partial rankings from multiple experts when only a limited number of relevant elements are explicitly ranked (top- K), and this number may vary from one expert to another (top- K_i). Moreover, we introduce two strategies to quantify the weight of each partial ranking. We cast this problem within the realm of combinatorial optimization and lean on the successful Ant Colony Optimization (ACO) metaheuristic algorithm to arrive at high-quality solutions. The proposed approach is evaluated through a real-world scenario and 190 synthetic datasets from www.PrefLib.org. The experimental evidence indicates that the proposed ACO-based solution is capable of significantly

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outperforming several evolutionary approaches that proved to be very effective when dealing with the Kemeny ranking problem.

Key words: Kemeny ranking problem, partial rankings, weighted aggregation, swarm intelligence, ant colony optimization.

1. Introduction

The aggregation of preferences from multiple experts is a well-studied topic in a number of fields such as *economic theory* (properties of a social choice function under elevation of pairs) [1], *social choice theory* (preference aggregation from a small subset of critical nodes in social networks) [2], *multi-criteria decision making* (group decision making) [3], *machine learning* [4] (evolutionary voting in classifier ensembles), *multi-agent systems* (reaching consensus in high-dimensional linear systems) [5] or *computational biology* (consensus genetic mapping) [6].

When these preferences are elicited in the form of N *rankings* over M objects/items (where each ranking denotes the preference of a single expert), the goal is to build a *consensus* (aggregated) ranking that reflects the set of individual preferences as faithfully as possible. Several methods to aggregate the ranking preferences of multiple voters have been proposed in the literature [7] [8] [9] [10]. Arrow’s axioms [11] state, however, that no aggregation method could simultaneously satisfy three fairness criteria: *non-dictatorship* (the voting results cannot simply mirror that of any single person’s preferences without consideration of the other voters), *Pareto efficiency* (if every individual prefers a certain option to another, then so must the resulting societal preference order) and *independence of irrelevant alternatives* (changes in individuals’ rankings of irrelevant alternatives –ones outside a certain subset– should have no impact on the societal ranking of the subset).

In spite of the above result, it is still possible to compute an *aggregated ranking* having minimal distance to the global consensus. This ranking is referred to as the *Kemeny ranking* [12] [13] and interpreted as a maximum likelihood estimator of the “correct” ranking. Unfortunately, the Kemeny ranking is NP-hard to calculate. The study in [14] thoroughly investigates different optimization methods (exact and approximate algorithms) for computing the Kemeny ranking. The authors concluded that heuristic approaches are recommended in contexts having weak or no consensus. More recently, Aledo et. al. [15] resorted to evolutionary algorithms to come to grips with this

challenging problem. Their results outperformed the remaining tested algorithms. Nevertheless, the proposed model is thought of for *complete rankings* (i.e., those in which each element is explicitly ranked) and cannot be directly applied to the aggregation of incomplete (*partial*) preferences, i.e., where only a subset of the available items is explicitly ranked. In this paper, we investigated two types of partial rankings that could be described as follows:

1. **Top- K rankings:** All respondents exactly select K relevant factors, whereas the remaining factors are placed at the $K+1$ position. In this kind of partial ranking, ties into the top- K ranked positions are not allowed.
2. **Top- K_i rankings:** Each respondent R_i is free to select K_i relevant factors that may be partially or completely ordered, whereas the remaining factors are placed at the K_i+1 position. In this scenario, tied factors into the top- K_i ranked positions could be observed.

This paper brings forth the following contributions. (1) We address the weighted aggregation of the two previous types of partial rankings from multiple experts by formulating the *weighted Kemeny ranking problem for partial rankings*, an extension of the Kemeny ranking problem that is able to aggregate top- K and top- K_i partial rankings from multiple experts. (2) We cast this problem into the realm of combinatorial optimization and lean on Ant Colony Optimization (ACO) [16], one of the most popular Swarm Intelligence [17] schemes, as the underlying optimization engine. In this scheme, we proposed two improved rules to compute the heuristic information used by ants to select the next state. (3) We introduce two heuristic strategies to derive the weight of each partial ranking in presence of subjective expert information (i.e., a set of predefined categories) or in its absence. In the first strategy, the weight is calculated from the fuzzy membership grade of each partial ranking to a set of predefined categories. If these predefined categories are not available, then the weight is computed as the ratio of non-tied items included in the partial ranking. (4) We conduct an extensive empirical analysis by comparing our solution against 11 other methods (two simple greedy techniques and 9 evolutionary optimizers) using a real-world scenario wherein Belgian respondents rank different aspects of potential employers, and 190 synthetic datasets taken from www.PrefLib.org. The empirical evidence indicates that the ACO-based approach is capable of significantly outperforming the other models for datasets under consideration.

The rest of the article is organized as follows. Section 2 briefly examines the Kemeny ranking problem and discusses several methods for aggregating partial rankings. Section 3 elaborates on a weighted extension of the Kemeny rule for aggregating partial rankings while Section 4 goes over the ACO fundamentals and revisits the three most prevalent models. Section 5 is concerned with tailoring ACO to solve the weighted Kemeny ranking problem, including the learning of the heuristic information matrix from the available data. Two heuristic strategies to compute the weight of each partial ranking are described in Section 6. The empirical study carried out to validate the proposed approach is unveiled in Section 7. Conclusions and future work directions are outlined in Section 8.

2. Related work and remaining challenges

This Section briefly reviews relevant works related to voting rules, the Kemeny ranking problem for complete rankings as well as other approaches for the aggregation of partial rankings.

2.1. Voting rules and Kemeny ranking problem for complete rankings

A *voting rule*, a.k.a *rank aggregation rule*, takes as input multiple rankings over the same element set and produces as outcome either a single element (the winner) or a consensus ranking of these elements [13].

Among the many different voting rules proposed in the literature [18], the *plurality rule* is perhaps one of the best known and most often applied *scoring rules*. This rule ranks items by the frequency with which they are placed first in the rankings. One may notice that other important considerations present in each ranking are simply disregarded by this procedure.

The *Borda rule* is another scoring rule. Each candidate earns as many points as the number of candidates ranked lower than himself. The winner is the one with the most points.

The *single transferable vote rule* goes through a series of $M - 1$ rounds, each one eliminating the element with the lowest plurality score from every ranking. The last remaining element is the winner.

The *Bucklin rule* computes a score for each element that is based on the number of voters that ranked it among the top- K candidates. An element “passes the post” if it is selected within the lowest K elements by at least half of the voters. Ties are broken by the number of votes by which the post is passed.

103 The *maximin rule* ranks elements after a score based on pairwise counts
 104 of the number of votes that placed that element higher than another element.

105 The *Copeland rule* also follows a score but this time an element earns/loses
 106 a point for every pairwise election it wins/loses.

107 The *ranked pairs rule* also returns a ranking based on an ordering of all
 108 element pairs (a,b) according to the number of voters that prefer a over b .

109 Another well-studied rule is the *Kemeny rule* [12] [15], which operates on
 110 complete rankings. This rule yields a ranking that maximizes the number of
 111 pairwise agreements among the individual rankings (votes), where a pairwise
 112 agreement is reached whenever the ranking agrees with one of the votes on
 113 which a pair of candidates is ranked higher [13]. More formally, given a set
 114 of N rankings $X = \{X_1, X_2, \dots, X_N\}$ over M elements, the *Kemeny ranking*
 115 *problem* is concerned with finding the ranking X_* that satisfies Equation (1),
 116 where \mathcal{P} stands for the set of all possible permutations over M elements
 117 (there are $M!$ possible permutations) and $\mathcal{K}(X_i, Y)$ denotes the Kendall-
 118 Tau distance between X_i and Y . The resultant ranking X_* is called the
 119 *Kemeny ranking* of the set and construed as the one minimizing the number
 120 of disagreements among all rankings in X [15].

$$X_* = \operatorname{argmin}_{Y \in \mathcal{P}} \frac{1}{N} \sum_{i=1}^N \mathcal{K}(X_i, Y) \quad (1)$$

121 2.2. Methods for aggregating partial rankings

122 González-Pachón and Romero [19] approach the aggregation of *quasi or-*
 123 *ders* (i.e., incomplete ordinal rankings) as a *consensus search* by using dis-
 124 tance functions. Interval goal programming (IGP) is presented as their solver
 125 of choice that tries to establish a weak consensus over incomplete ordinal
 126 rankings.

127 Klementiev et. al. [20] proposed a rank aggregation method for both
 128 permutations and top- K lists where they account for the *type* of the ele-
 129 ments being ranked, i.e., they could belong to different data domains, so as
 130 to include the notion of *domain expertise*. Given only a set of constituent
 131 rankings, they learn an aggregation function that attempts to recreate the
 132 true ranking without labeled (type) data. The method is based on a mixture
 133 of distance-based models and leans on the Expectation-Maximization (EM)
 134 algorithm to estimate its parameters. The new technique significantly and
 135 robustly outperformed their previous domain-agnostic model [21].

136 Ammar and Shah [22] consider partial data in the form of *first-order*
 137 or *comparison* marginals. They treat this information as partial samples
 138 from an unknown distribution over permutations and provide an efficient
 139 algorithm for finding an aggregate complete ranking directly from the data
 140 without first learning the underlying distribution; this is an appealing feature
 141 for designing large-scale ranking systems such as *recommendation systems*.

142 Neghaban et. al. [23] remarked that the approach in [22] requires in-
 143 formation about comparisons between all element pairs, and for each pair
 144 it requires the exact pairwise comparison marginal w.r.t the underlying per-
 145 mutation distribution. This assumption is not always easy to meet since,
 146 in reality, all pairs of items are not usually compared. The authors then
 147 propose an algorithm that takes as input the noisy comparison marginals for
 148 a subset of all possible item pairs and spits out scores for each item. The
 149 noise in the underlying permutation distribution is modeled after the Multi-
 150 nomial Logit (MLN) method [24]. Their algorithm has a natural *random*
 151 *walk* interpretation over the graph of objects with edges present between
 152 two objects if they are compared; the scores turn out to be the stationary
 153 probability of this random walk. The empirical analysis indicates that the
 154 proposed scheme performs comparably to the maximum likelihood estimator
 155 of the MLN model and outperforms the technique in [22].

156 Brandenburg et. al. [25] studied the aggregation of partial rankings
 157 under the nearest neighbor (NN) and Hausdorff versions of the Kendall-Tau
 158 distance. They proved that this problem is **NP-complete** under the NN
 159 Kendall-Tau distance even for two voters and that, in contrast, it is **NP-**
 160 **hard** and **coNP-hard** under the Hausdorff Kendall-Tau distance for at least
 161 four voters.

162 2.3. Remaining challenges

163 In spite of the Arrow’s impossibility theorem, researchers continue ad-
 164 dressing the aggregation of several preferences by solving the Kemeny rank-
 165 ing problem. Young and Levenglick [26] show that the Kendall distance (and
 166 consequently its extensions) is the only distance function ensuring the per-
 167 mutation(s) minimizing the Kemeny ranking problem have three desirable
 168 properties of being *neutral*, *consistent* and *Condorcet*. The Condorcet prop-
 169 erty means that, if there exists a permutation such that the order of every
 170 pair of elements is the order preferred by the majority, then that permu-
 171 tation has minimum distance to the voters’ permutations. Therefore, the

172 main challenge towards this goal lies on the performance of the discrete op-
 173 timizer used when solving the related combinatorial problem. On the other
 174 hand, there exist situations for which rankings are partial and therefore, the
 175 classical Kendall-Tau distance is no longer suitable.

176 The second challenge refers to the inclusion of the weighted approach
 177 when aggregating partial rankings and the automatic estimation of the mem-
 178 bership degree of a ranking to the population. Recently Nápoles et al. [27]
 179 proposed a two-step methodology to build fuzzy prototypes from a popula-
 180 tion of partial rankings. Being more explicit, in the first step the authors
 181 put forth a fuzzy clustering algorithm for partial rankings called fuzzy c -
 182 aggregation, while the second step is focused on solving the extended Ke-
 183 meny ranking problem for each discovered cluster taking into account the
 184 estimated partition matrix. Despite the novelty of this approach, the reader
 185 may notice that this algorithm will produce c different aggregations, with c
 186 being the number of clusters detected by the clustering algorithm. However,
 187 the clustering approach may not be adequate for some scenarios where a sin-
 188 gle aggregated solution is expected. This implies that other approaches to
 189 compute the membership degrees of partial rankings are required.

190 3. Weighted aggregation of partial rankings

191 In this section we extend the well-known Kemeny ranking problem [12]
 192 by considering that orderings to be clustered may be incomplete or partial
 193 (i.e., tied elements are allowed). Besides, we assume that each partial or-
 194 dering X_i has an associated weight $\omega_i \in [0, 1]$ representing the extent to
 195 which the ranking belongs to the population. Formally, *the weighted Ke-*
 196 *meny ranking problem for partial rankings* could be summarized as follows.
 197 Let $X = \{X_1, \dots, X_i, \dots, X_N\}$ be a set of N partial rankings over M items
 198 $F = \{F_1, \dots, F_l, \dots, F_M\}$ where the i th ranking comprises the vote of a single
 199 respondent with weight $\omega_i \in [0, 1]$. More explicitly, we can describe a partial
 200 ranking X_i as a vector $\{X_i^1, \dots, X_i^k, \dots, X_i^M\}$ where $X_i^k \preceq X_i^{k+1}$ denotes that
 201 X_i^k precedes X_i^{k+1} . The theoretical challenge is to construct a fair enough
 202 ranking Y taking into account all input (potentially partial) rankings and
 203 their weights. It should be mentioned that the solution for the weighted
 204 Kemeny ranking problem is a complete ranking, and therefore ties are not
 205 allowed.

206 Being more explicit, the solution for the weighted Kemeny problem is
 207 equivalent to computing a complete ranking with minimal distance to the

208 global consensus. Equation (2) formalizes the objective function to be min-
 209 imized, where $\mathcal{H}(X_i, Y)$ represents a distance function quantifying the dis-
 210 similarity between the i th partial ranking and the candidate solution Y to
 211 be evaluated.

$$\min \rightarrow \mathcal{F}(Y) = \sum_{X_i \in X} \omega_i \mathcal{H}(X_i, Y) / \sum_i \omega_i \quad (2)$$

212 The reader may notice that the first modification to the standard Kemeny
 213 ranking problem lies on the inclusion of the weight quantifying the extent to
 214 which the i th ranking belongs to the population. The second modification
 215 is related to the normalized distance function $\mathcal{H}(\cdot, \cdot)$ to compute the dis-
 216 similarity degree between two rankings with tied elements. Notice that the
 217 standard Kemeny ranking problem uses the Kendall-Tau distance [28], which
 218 measures the dissimilarity as the number of item pairs over which the two
 219 rankings disagree. However, the original Kendall-Tau distance is no longer
 220 adequate when comparing rankings having tied items since this distance as-
 221 sumes that items are all ordered. Instead, we could adopt other versions
 222 of the Kendall-Tau distance or other extended dissimilarity measures such
 223 as the Hausdorff distance [29], the Spearman's footrule distance [30] or the
 224 Goodman-Kruskal's one [31]. Having several metrics for partial rankings is
 225 obviously convenient, but it poses the question of which one would be better
 226 suited when comparing partial rankings when solving the Kemeny ranking
 227 problem.

228 The Goodman-Kruskal's approach is not always defined and thus there
 229 could be scenarios where this procedure fails. Moreover, Fagin et al. [32]
 230 mathematically proved that the Hausdorff variants of the Kendall-Tau dis-
 231 tance and the Spearman's footrule distance are actually equivalent. This out-
 232 come was based on the Diaconis-Graham inequality [33], which asserts that
 233 the Kendall-Tau distance and the Spearman's footrule distance are within a
 234 factor of two from each other. It implies that selecting a distance function
 235 does not matter so much when solving the weighted Kemeny ranking prob-
 236 lem, as long the distance function is capable to deal with partial rankings.

237 In this paper we use the Hausdorff version of the Kendall-Tau distance
 238 as the dissimilarity functional when aggregating partial rankings. The Haus-
 239 dorff distance has been extensively studied and shown to have particularly
 240 flexible mathematical and algorithmic properties [32]. Equation (3) formal-
 241 izes this distance, where X_i is the i th partial ranking, Y denotes the Kemeny

242 ranking to be evaluated, $\mathcal{K}(X_i, Y)$ is the set of all item pairs that appear in
 243 different order, $\mathcal{R}_1(X_i, Y)$ is the set of all item pairs which are tied in X_i
 244 but not tied in Y , while $\mathcal{R}_2(X_i, Y)$ is the set of all item pairs which are tied
 245 in the ranking Y but not tied in the i th partial ranking. This function can
 246 also be adopted for comparing full rankings, thus leading to the Kendall-Tau
 247 distance (i.e., $\mathcal{R}_1(X_i, Y) = \mathcal{R}_2(X_i, Y) = 0$).

$$\mathcal{H}(X_i, Y) = |\mathcal{K}(X_i, Y)| + \max \{|\mathcal{R}_1(X_i, Y)|, |\mathcal{R}_2(X_i, Y)|\} \quad (3)$$

248 The inclusion of the Hausdorff distance $\mathcal{H}(X_i, Y)$ in the objective function
 249 (2) allows computing the dissimilarity between each input ranking and the
 250 candidate (complete) aggregated ranking. Due to the fact that Y is a full
 251 ranking, we could compute $\mathcal{H}(X_i, Y) = |\mathcal{K}(X_i, Y)| + |\mathcal{R}_1(X_i, Y)|$. This is
 252 possible because there are no tied items in a full ranking (i.e., $|\mathcal{R}_2(X_i, Y)| =$
 253 0). On the other hand, the reader may verify that $|\mathcal{R}_1(X_i, Y)| = \binom{M-K_i}{2}$
 254 where K_i is the number of relevant items selected by the i th respondent.
 255 Notice that we assume partial rankings with non-homogeneous tied factors,
 256 since there are scenarios where each respondent may select a different number
 257 of relevant items. Equation (4) shows the normalized objective function to
 258 be optimized.

$$\min \rightarrow \mathcal{F}(Y) = \left(\sum_{X_i \in X} \frac{2\omega_i [|\mathcal{K}(X_i, Y)| + \binom{M-K_i}{2}]}{M(M-1)} \right) / \sum_i \omega_i \quad (4)$$

259 Equation (4) involves a **NP-hard** problem with a search space comprised
 260 of $M!$ possible states (i.e., the set of all permutations over M items). In or-
 261 der to deal with the computational intractability of this weighted aggregation
 262 problem, Nápoles et al. [34] proposed a novel approach based on Swarm Intel-
 263 ligence that exploits a colony of artificial ants. However, this approach does
 264 not take into account the weight of partial rankings. Recently, Nápoles et al.
 265 [27] extended the crisp method in order to construct prototypes from fuzzy
 266 information granules discovered by a clustering algorithm. Before describing
 267 the details of this procedure, next we provide a basic background about Ant
 268 Colony Optimization that will be used to solve the weighted Kemeny ranking
 269 problem formulated before.

270 4. Ant Colony Optimization

271 The generation of feasible permutations representing complete rankings
 272 is entrusted in this study to the ACO methods. The objective function
 273 in Equation (4) evaluates the quality of each candidate solution (ant tour)
 274 during the search process.

275 The ACO metaheuristic is a biologically-inspired search technique that
 276 was originally devised to solve combinatorial optimization problems [16]. Its
 277 creator, Marco Dorigo, drew inspiration from the manner in which ants cor-
 278 porately forage. They depart from the nest and once a source of food is iden-
 279 tified, they deposit a chemical substance on the ground named *pheromone*
 280 on their way back to the nest; these pheromone trails serve to guide the
 281 rest of the colony towards the food source [35]. ACO is one of the hallmark
 282 swarm intelligence algorithms and bears a plethora of successful applications
 283 to real-world problems [36] [37] [38].

284 ACO is a fully constructive model where each ant builds a candidate so-
 285 lution to the problem by incrementally exploring the nodes (or edges) of a
 286 search graph. Each artificial ant moves from one state to another during
 287 the search process (here states are components of the solution). As depicted
 288 in Equation (5), the likelihood of moving from one node to another (ACO
 289 transition rule) at the next discrete time step $t+1$ mainly rests on two param-
 290 eters: (1) the *collective information* $\tau_{kl}(t)$ derived from the pheromone trails
 291 and iteratively updated by ants during the navigation of the search graph
 292 and (2) the *heuristic information* η_{kl} denoting the invariant, problem-specific
 293 preference of moving from one state to another. The heuristic component
 294 must be carefully provided/estimated as it is treated as an invariant, i.e., it
 295 is not modified throughout the algorithm’s execution.

$$P_{kl}^v(t+1) = \frac{[\tau_{kl}(t)]^\alpha [\eta_{kl}]^\beta}{\sum_{r \in \mathcal{N}_k^v} [\tau_{kr}(t)]^\alpha [\eta_{kr}]^\beta}, \quad l \in \mathcal{N}_k^v \quad (5)$$

296 In light of the Kemeny ranking problem, Equation (5) denotes the prob-
 297 ability of accepting the l th state (i.e., next factor to be ranked) at the k th
 298 position of the candidate ranking, \mathcal{N}_k^v is the set of unvisited states (factors)
 299 at the k th position for the v th ant while α and β govern the strength of the
 300 collective and heuristic information, respectively.

301 Once the individual ant tours are completed, the pheromone levels on
 302 all trails using the solutions found by the agents will be updated. First,

303 pheromone evaporation takes place uniformly thus reducing the amount of
 304 pheromone on all trails. Subsequently, certain pheromone amount will be
 305 added to the nodes/edges of the more promising solution(s). This is a very
 306 important step in any ACO-based implementation; most of ACO variants
 307 differ mainly in the strategy used for updating the collective information
 308 (pheromone trails) at each iteration. In the sequel we will discuss the three
 309 most popular ACO algorithms.

310 4.1. *Ant System*

311 *Ant System* (AS) is credited with being the first ACO algorithm [39].
 312 The pheromone trails are updated once all ants have completed their tours.
 313 A certain portion of the pheromone in each trail is evaporated according to
 314 a factor $0 < \rho < 1$. Afterwards, each ant v deposits a pheromone amount
 315 $\Delta\tau_{kl}^v$ proportional to the quality of its solution along the edges belonging to
 316 it. This pheromone update rule is reflected in Equation (6), with S being
 317 the number of ants in the colony.

$$\tau_{kl}(t+1) = (1 - \rho)\tau_{kl}(t) + \sum_{v=1}^S \Delta\tau_{kl}^v \quad (6)$$

318 The long-term effect of the above rule is that edges not frequently chosen
 319 by the ants will see their pheromone concentration gradually vanish whereas
 320 those edges selected by the ants will receive a boost in their pheromone
 321 amount, thus becoming more probable candidates for selection in future it-
 322 erations. A more thorough study [39] revealed that better results could be
 323 attained if the pheromone increase is only applied by the global best solution
 324 rather than having all colony members do so. In spite of that, AS suffers from
 325 stagnation (convergence to local optima) due to the unbounded accumulation
 326 of pheromone over the best found edges.

327 4.2. *Ant Colony System*

328 *Ant Colony System* (ACS) improves the AS scheme by exploiting the
 329 global best solutions found during the search stage [40]. As a result, the
 330 algorithm exhibits superior exploitation features as ants build their solutions,
 331 instead of exploring new areas of the solution space. This goal is achieved via
 332 a three-fold mechanism: (1) a strong *elitist* strategy for updating pheromone
 333 trails, (2) a modification to the pheromone update rule and (3) a pseudo-
 334 random rule for selecting new states.

335 ACS' pheromone update rule is reported in Equation (7), with $\tau_{kl}^*(t)$
 336 denoting the pheromone amount associated with the ant featuring the best
 337 heuristic value at time step t . Like in AS, pheromone evaporation affects all
 338 edges yet the boost is only reserved for those edges belonging to the best
 339 solution.

$$\tau_{kl}(t+1) = (1 - \rho)\tau_{kl}(t) + \rho\tau_{kl}^*(t) \quad (7)$$

340 ACS' pseudo-random proportional rule in Equation (8) aims at fostering
 341 exploitation of the knowledge attained by the colony. In a nutshell, if a
 342 random number $q \sim U(0, 1)$ falls below q_0 then the ant will move to the
 343 state maximizing the product between collective and heuristic information,
 344 otherwise ACS will adopt the standard decision rule in Equation (5). Notice
 345 that q_0 is a user-defined parameter that favors exploitation over exploration
 346 as it approaches 1.

$$l = \operatorname{argmax}_{r \in \mathcal{N}_k^v} \{[\tau_{kl}(t)]^\alpha [\eta_{kl}]^\beta\} \text{ if } q \leq q_0 \quad (8)$$

347 The third distinctive element in the ACS model is the iterative pheromone
 348 update rule ants employ as they build their solution, as shown in Equa-
 349 tion (9). This approach has the same effect as decreasing the probability
 350 of selecting the same path for all ants, thus introducing a balance between
 351 exploitation and exploration.

$$\tau_{kl}(t+1) = (1 - \rho)\tau_{kl}(t) + \rho\tau_{kl}^*(0) \quad (9)$$

352 The ACS algorithm frequently reports better performance than AS owing
 353 to its emphasis on the exploitation of the most promising solutions discovered
 354 by the colony.

355 4.3. MAX-MIN Ant System

356 Like ACS, the *MAX-MIN Ant System* (MMAS) [41] was specifically en-
 357 gineered to pursue a stronger exploitation of solutions, thus avoiding the
 358 stagnation problems encountered by AS. This model has the following fea-
 359 tures: (1) similar to ACS, a strong elitist strategy regulates the ant allowed
 360 to update the pheromone trails (either the best-so-far ant or the one with the
 361 best solution in the current iteration); (2) all pheromone trails are bound to
 362 the range $[\tau_{MIN}, \tau_{MAX}]$. If $\tau_{MIN} > 0$ for all solution components, then the

probability of choosing a specific state will never be zero, which avoids stagnation configurations. Finally, pheromone trails are initialized with τ_{MAX} to ensure further exploration of the search space at the beginning of the optimization phase.

MMAS has also reported very encouraging results in the literature, even outperforming ACS [42] [36].

5. Solving the weighted Kemeny ranking problem

In this section we explain how to optimize the objective function defined in Section 3 by exploiting a colony of artificial ants. With this goal in mind, we defined four central components:

- The structure of the pheromone graph used by ants to construct the solutions.
- The interpretation of the probabilistic rule to select the next state.
- The formal definition of the set of feasible states at each step.
- The estimation of the heuristic information.

As mentioned, the goal of the search method is to produce a complete ranking (i.e., a permutation over M different factors) minimizing the objective function (4). This problem is similar to the well-known Traveling Salesman Problem [43] where artificial agents construct the candidate solution by traveling along a fully connected graph. The graph nodes correspond to the M elements $F = \{F_1, ..., F_l, ..., F_M\}$ to be ordered. Due to the fact that a solution to the weighted Kemeny ranking problem is a permutation of such M items, each item F_l will appear exactly once. This suggests that self-connected graph nodes are not allowed, otherwise the Kemeny ranking may involve explicit tied items. However, a Kemeny ranking might comprise implicit tied items (i.e., items that may be freely exchanged without altering the heuristic value) which is a result of frequent ties over the same two items.

In the proposed scheme, the transition value P_{kl} is the probability of accepting the l th item at the k th ranking position. This approach is slightly different from other scenarios where the transition value P_{kl} often denotes the probability of moving to the l th graph node from the k th node. In practice, both approaches are equivalent because the probability of accepting

the l th ranking item at the k th position will eventually be influenced by those ranking items situated at the previous $(k-1)$ positions. From this remark we can formally define the set of feasible states for the v th ant at each step k . The domain set $\mathcal{N}_k^v \subseteq F$ is given by $\mathcal{N}_k^v = \{F_1, \dots, F_l, \dots, F_M\} - \{Y_v^1, Y_v^2, \dots, Y_v^{k-1}\}$ where Y_v^{k-1} represents the item located at the $(k-1)$ ranking position, according to the v th agent. Being more explicit, all previously ranked items are no longer part of the neighborhood of the ant at the k th step. This ensures the unicity of ranking items in the solution for the Kemeny ranking problem.

Another important aspect when solving combinatorial problems using ACO-based algorithms is the estimation of the heuristic matrix. The accurate estimation of the heuristic component often leads to high-quality solutions, otherwise the solutions to the weighted Kemeny ranking problem will probably be sub-optimal. In the following sections we propose two strategies to estimate the heuristic matrix from input data, assuming two partial ranking aggregation scenarios.

5.1. *Weighted aggregation of multiple top- K rankings*

The first scenario takes place when each respondent selects the top- K items. It implies that each input ranking will be partial in the sense that only the top- K items are explicitly ranked, whereas the other $M-K$ items are tied at the $K+1$ position. It can be noticed that estimating the heuristic values for the M items across the first K positions is equivalent to computing the number of observations on which the l th element *was observed* at the k th position ($k = 1, 2, \dots, K$). For the $M-K$ last positions, this heuristic cannot be directly used since such items are tied. However, we may compute the number of observations on which the item *was not included* into the top- K .

Equation (10) formalizes the above reasoning, where $\vartheta_k(F_l)$ denotes the sum of the weights of those rankings on which item F_l was ranked at the k th position ($1 \leq k \leq K$), while $\sim \vartheta_k(F_l)$ is the sum of the weights of those rankings on which the l th item was not included into the top- K . For the latter case divide the expression by $M-K$ since these items have probability $(M-K)/M$ to be placed at the last positions, assuming a uniform probability distribution.

$$\eta_{kl} = \begin{cases} \vartheta_k(F_l) \left(\sum_i \omega_i \right)^{-1}, & k \leq K \\ \frac{\sim \vartheta_k(F_l)}{M - K} \left(\sum_i \omega_i \right)^{-1}, & k > K \end{cases} \quad (10)$$

428 It should be specified that the functions $\vartheta_k(F_l)$ and $\sim \vartheta_k(F_l)$ must con-
 429 sider the fact that X_i belongs to the ranking population with weight ω_i .
 430 Equation (11) show how to compute the function $\vartheta_k(F_l)$, but it may be eas-
 431 ily extended for $\sim \vartheta_k(F_l)$. On the other hand, in Equation (10) the sum of
 432 all weights $\sum_i \omega_i \leq N$ is used to normalize the heuristic values.

$$\vartheta_k(F_l) = \sum_{X_i \in X} \begin{cases} \omega_i & , X_i^k = F_l \\ 0 & , X_i^k \neq F_l \end{cases} \quad (11)$$

433 **Example 1.** Let us consider a weighted aggregation scenario of $N = 5$
 434 partial rankings over $M = 5$ items where each expert selected the $K = 3$ most
 435 relevant factors. Table 1 summarizes this scenario, where each row involves
 436 a partial ranking. According to Equation (10), the heuristic preference of
 437 accepting the item F_2 at the second position is given by $\eta_{22} = \vartheta_2(F_2)/2.2 =$
 438 $0.8/2.2 \approx 0.36$. Similarly we can compute the remaining components of the
 439 heuristic matrix.

Table 1: Example of a weighted aggregation of multiple top- K rankings.

	ω_i	F_1	F_2	F_3	F_4	F_5
X_1	0.8	1	2	$K + 1$	3	$K + 1$
X_2	0.2	2	1	$K + 1$	$K + 1$	3
X_3	0.5	1	3	$K + 1$	$K + 1$	2
X_4	0.1	1	3	$K + 1$	2	$K + 1$
X_5	0.6	2	1	3	$K + 1$	$K + 1$

440 Observe that according to Equation (10) some heuristic values could be
 441 zero (e.g. $\eta_{25} = 0$ since F_2 was always included into the top-3) and therefore
 442 the probability P_{kl}^v of selecting these states will be zero. However, normally
 443 the probability value P_{kl}^v should not be zero since it is possible to estimate a
 444 good solution having F_2 at the last position. In order to overcome this issue
 445 we replace all zero-values by $\eta_{MIN} = \min \{\eta_{kl}\}$ such that $\eta_{kl} \neq 0$, thus we

446 guarantee that all states have a probability to be visited by artificial ants,
 447 although they have less chance to be selected.

448 5.2. Weighted aggregation of multiple top- K_i rankings

449 This scenario is more complex (but also more informative) because each
 450 respondent is free to select K_i items such that $2 \leq K_i \leq M$. Since the number
 451 of relevant items could change from a respondent to another, we cannot
 452 simply count the number of observations of each item. In order to compute
 453 a more realistic heuristic matrix we may compute the relative frequency on
 454 which an item could be observed at each ranking position. Notice that this
 455 assumption attempts to hypothetically break the ties in order to transform
 456 partial rankings into complete ones. Equation (12) enunciates this method,
 457 where $Q_{K_i}(F_l)$ is the set of all rankings where item F_l was excluded from
 458 the top- K_i , $\psi_{X_i}(F_l)$ is the set of all feasible positions for the l th item, while
 459 $\vartheta_k(F_l)$ denotes the sum of the weights of those rankings on which item F_l
 460 was placed at the k th ranking position.

$$\eta_{kl} = \left(\vartheta_k(F_l) + \sum_{X_i \in Q_{K_i}(F_l)} \begin{cases} \omega_i & , k \in \psi_{X_i}(F_l) \\ 0 & , k \notin \psi_{X_i}(F_l) \end{cases} \right) \left(\sum_i \omega_i \right)^{-1} \quad (12)$$

461 **Example 2.** Let us consider the weighted aggregation scenario summa-
 462 rized in Table 2, with $N = 5$ partial rankings over $M = 5$ items where the
 463 i th respondent selected the most relevant K_i items. According to Equa-
 464 tion (12), the heuristic value of accepting F_2 at the first position η_{12} is
 465 $1/2.2(0.7) \approx 0.318$ since $\vartheta_1(F_2) = 0.7$, $Q_{K_i}(F_2) = \{X_1, X_4\}$, $\psi_{X_1}(F_2) =$
 466 $\{4, 5\}$, $\psi_{X_4}(F_2) = 3, 4, 5$. Observe that item F_2 could not be hypothetically
 467 located at the first ranking position without introducing new tied pairs of
 468 items because $k \notin \psi_{X_1}(F_2) \cup \psi_{X_4}(F_2)$. This suggests that the heuristic value
 469 η_{12} is computed from the number of times F_2 was observed at the first posi-
 470 tion.

471 Similarly to the first scenario (i.e., respondents select the most relevant K
 472 items), we must avoid zero-values in the heuristic matrix, although this situa-
 473 tion is possible (i.e., the item was never observed in a position and there is no
 474 chance to be observed without inducing new ties). However, it is still possible
 475 to build a candidate solution with this feature having minimal distance to
 476 the consensus, and therefore it must be considered as well. In these scenarios
 477 the probability should not be zero but rather small, e.g., $\eta_{MIN} = \min\{\eta_{kl}\}$

Table 2: Example of a weighted aggregation of multiple top- K_i rankings.

	ω_i	F_1	F_2	F_3	F_4	F_5
X_1	0.8	1	$K_1 + 1$	2	$K_1 + 1$	3
X_2	0.2	3	1	$K_2 + 1$	2	$K_2 + 1$
X_3	0.5	2	1	4	5	3
X_4	0.1	2	$K_4 + 1$	$K_4 + 1$	$K_4 + 1$	1
X_5	0.6	5	3	1	2	4

where $\eta_{kl} \neq 0$. This approach is similar to the thresholding strategy used in MAX-MIN Ant Systems [41] which proved to be quite effective in promoting the exploration of alternative regions of the search space. Next we propose two strategies to compute the weight of each partial ranking.

6. Heuristics to determine the weight of each partial ranking

A pivotal issue when solving the weighted aggregation problem is related to the estimation of weights. Strategies for weighting partial rankings could vary from expert-based estimations to more automated measures. In this paper we addressed this issue by considering two scenarios. In the first one, the weight is calculated from the fuzzy membership grade of each partial ranking to a set of predefined categories. If these predefined categories are not available (second scenario), then the weight is computed as the ratio of non-tied items included in the partial ranking. This latter heuristic is based on the fact that partial rankings having a fewer number of tied factors are more informative when solving the weighted aggregation problem, as they better express the user preferences. Note however that we are assuming that partial ranking instances have the same confidence level (i.e., all experts responses are equally reliable).

Next we describe an algorithm to compute the membership degree of each partial top- K or top- K_i ranking across a set of predefined categories. In this paper, we assume that a category is a disjoint set of items that comprises an information granule regularly defined by domain experts. This *fuzzy allocation problem* could be formalized as follows. Let us suppose a set of N partial rankings $X = \{X_1, \dots, X_i, \dots, X_N\}$ over M items $F = \{F_1, \dots, F_l, \dots, F_M\}$ and a set of categories $C = \{C_1, \dots, C_j, \dots, C_P\}$ resulting from a partition of the item set. The fuzzy allocation problem is equivalent to compute the degree on which partial rankings belong to each predetermined category. This allo-

505 cation problem is fuzzy in nature because a partial ranking may be associated
 506 with several categories at the same time but with different degrees.

507 Essentially, the proposed method computes the correspondence degree
 508 between items in the partial ranking and those items belonging to each cat-
 509 egory. Observe that we cannot compute this correspondence degree using a
 510 distance function (e.g., the Hausdorff distance) since categories are unordered
 511 sets and therefore there is no ordinal relation among category items. The
 512 method comprises four well-defined steps which are described below.

513 **Step 1.** Compute the intersection set $\Phi_{ij} = X_i \cap C_j$ between the partial
 514 ranking X_i and the category C_j . This step allows determining, for each
 515 partial ranking, the set of ranked items that additionally belongs to the j th
 516 category. Notice that this step does not consider the existence of an ordinal
 517 relation between pairs of items included into the top- K (or top- K_i).

518 **Step 2.** Compute the relative relevance $Z(F_l)$ of each factor $F_l \in \Phi_{ij}$
 519 using its position $1 \leq R_{(F_l)} \leq K$ into the top- K_i (or top- K) items associated
 520 with the i th partial ranking. The relevance $Z_i(F_l) = [(K_i + 1) - R_{(F_l)}]/K_i$
 521 provides a local measure to determine the degree of membership to each
 522 category. In the case of top- K partial rankings, $K_1 = \dots = K_i = \dots = K_N$
 523 since all respondents have to select exactly K relevant items.

524 **Step 3.** Compute the weight $\tilde{\omega}_i^{(j)}$ of the i th partial ranking to the j th
 525 category. To accomplish that, we adopt Equation (13) for both top- K and
 526 top- K_i scenario, assuming that ψ is the normalization factor.

$$\tilde{\omega}_i^{(j)} = \sum_{F_l \in \Phi_{ij}} \frac{Z(F_l)}{\psi} \quad (13)$$

527 It should be remarked that, for top- K partial rankings, the number of
 528 relevant items K may be different from the cardinality of the category C_j . If
 529 $|C_j| < K$ then the degree to which the i th partial ranking belongs to the j th
 530 category will never be maximal because some ranking positions cannot be
 531 covered. On the other hand, if $|C_j| > K$ the degree to which the i th partial
 532 ranking belongs to the j th category will never be maximal either because
 533 some ranking items cannot be selected by respondents. Both scenarios are
 534 considered when normalizing the sum of all relevance degrees, which allows
 535 computing fair membership degrees.

536 Therefore, for top- K partial rankings, $\psi = \frac{-\min\{K, |C_j|\}[-2K + \min\{K, |C_j|\} - 1]}{2K}$.
 537 This normalization factor represents the sum of the first $\min\{K, |C_j|\}$ rele-
 538 vance degrees. Being more explicit, this sum of relevance degrees is equivalent

539 to computing the sum of the first K numbers $i/K, i = \{1, \dots, K\}$ minus the
 540 sum of the $K - \min\{K, |C_j|\}$ numbers $i/K, i = \{1, \dots, K\}$. It implies that
 541 the maximal value for the sum of the relevance degrees $Z(F_l), \forall F_l \in \Phi_{ij}$ is
 542 reached when $C_j \subseteq X_i$ and items contained in C_j are placed at the first K
 543 ranking positions. The normalization factor allows estimating realistic values
 544 and may be inferred from the following expression:

$$\begin{aligned}
 & \left(\sum_{i=1}^K i/K \right) - \left(\sum_{i=1}^{K-\min\{K, |C_j|\}} i/K \right) \\
 &= \frac{K+1}{2} - \frac{(K - \min\{K, |C_j|\})(K - \min\{K, |C_j|\} + 1)}{2K} \\
 &= \frac{-\min\{K, |C_j|\}[-2K + \min\{K, |C_j|\} - 1]}{2K}
 \end{aligned}$$

545 In the case of top- K_i partial rankings, the number of relevant items will
 546 likely differ from one partial ranking to another. More importantly, some
 547 of these items may be placed at the same ranking position (i.e., tied items
 548 into the top- K_i are allowed). This feature increases the uncertainty in the
 549 decision-making process and may lead to quite similar membership degrees.
 550 If the relevant items are uniformly selected from homogenous categories,
 551 then all membership degrees will probably tend to the membership value
 552 $1/|C|$. This configuration cannot be observed in the previous scenario since
 553 we assumed that the top- K items are rigorously ordered.

554 Another issue that arises here is that there is no restriction on the number
 555 of relevant items to be selected by the respondent when constructing the
 556 top- K_i partial ranking. This implies that a specific category can be entirely
 557 included into the top- K_i ranking. It could be possible to allocate all selected
 558 items at the same ranking level (e.g., selected items belong to the same
 559 category and they are equally relevant to characterize the concept under
 560 evaluation). In this case, the normalization factor $\psi = |C_j|$.

561 **Step 4.** After computing the above equation for each category, the final
 562 membership values are calculated as $\omega_i^{(j)} = \tilde{\omega}_i^{(j)} / \sum_j \tilde{\omega}_i^{(j)}$. This ensures that
 563 the sum of all membership values will be exactly one, which is an important
 564 property to be preserved in fuzzy environments.

565 7. Numerical simulations

566 In this section, we evaluate the proposed weighted aggregation approach
 567 across several evolutionary approaches that proved to be adequate solvers
 568 of the Kemeny ranking problem. With this goal in mind, we used both
 569 real-world and synthetic datasets having different features.

570 7.1. Benchmarking algorithms and parametric settings

571 In this section, we describe the algorithms selected for benchmarking pur-
 572 poses. Recently Aledo et al. [15] proposed a solution scheme based on Evolu-
 573 tionary Computation for the Kemeny ranking problem. Results have shown
 574 that evolutionary algorithms clearly outperformed other tested algorithms
 575 (i.e., Borda counting index, variants of the Branch and Bound algorithm,
 576 among others). The central feature of the Genetic Algorithms (GA) used to
 577 solve the Kemeny ranking problem relies on the search space characteristics.
 578 Instead of dealing with the standard binary representation, they adopted a
 579 permutation-based solution encoding. During the search progress, the chro-
 580 mosome population evolves according to three genetic operators: selection,
 581 crossover and mutation. The selection operator promotes high-quality in-
 582 dividuals and does not depend on the solution representation, but on the
 583 fitness value.

584 Nevertheless, in permutation-based search spaces crossover and mutation
 585 operators must be carefully defined; otherwise non-feasible solutions may be
 586 produced. In this study, we used the operators discussed in [44] to solve the
 587 *Traveling Salesman Problem*. Once promising individuals have been selected,
 588 they are (randomly) organized in pairs. Over each pair we apply the crossover
 589 operation by using one of the following three operators:

- 590 • **POS.** Position-based crossover operator [45]. It starts by selecting a
 591 set of random positions such that the values for these positions are kept
 592 in both parents. The remaining positions are completed by using the
 593 relative order in the other parent.
- 594 • **OX1.** Order crossover operator [46]. It selects two cut points $1 \leq$
 595 $c_1 < c_2 \leq M$ and then, for every parent, the genetic segment between
 596 the cut points c_1 and c_2 is directly copied into the corresponding child.
 597 Then, starting from c_2 , the remaining items are copied into that child
 598 following the relative order in which they appear in the other parent,
 599 onwards and moving to the first position once the end of the individual

600 is reached. The same procedure is repeated in the other child, by
 601 exchanging the role between the parent and the child.

- 602 • **OX2.** Order-based crossover operator [45]. It randomly selects several
 603 positions for each parent. Items in non-selected positions are main-
 604 tained in the child, while the selected ones are set according to the
 605 positions taken by these items in the other parent.

606 Once offspring are generated by crossover, a mutation operator is applied
 607 over each offspring with a given mutation probability. In [13] the authors
 608 adopted the following evolutionary operators:

- 609 • **ISM.** Insertion mutation operator [47]. It randomly chooses an element
 610 in the permutation, which is removed and reinserted in a different (ran-
 611 domly selected) position.
- 612 • **DM.** Displacement mutation operator [48]. It randomly selects a seg-
 613 ment of items in the permutation, which are removed and reinserted in
 614 a randomly selected position.
- 615 • **IVM.** Inversion mutation operator [49]. The semantic of this genetic
 616 operator is quite similar to DM but the removed items are reinserted
 617 as a single block in reverse order.

618 The combinations of the above permutation-based crossover and muta-
 619 tion operators lead to nine GA-based optimizers. For such evolutionary ap-
 620 proaches, we used a population of 200 individuals. The mutation probability
 621 is set to 0.1 whereas the crossover probability was fixed to 0.9. Observe that
 622 the crossover probability is notably higher than the mutation probability as
 623 suggested in [15]. The search process stops after 50 generations, leading to
 624 10,000 evaluations of the objective function. Normally, the number of gen-
 625 erations used in population-based algorithms is higher than the number of
 626 individuals. However, after a number of preliminary simulations we observed
 627 that, for the same number of generations, the GA-based algorithms reported
 628 better results by using a larger population and fewer generations.

629 In the case of ACO-based algorithms, we adopted the following param-
 630 eters: the number of ants was taken as the number of items to be ranked
 631 multiplied by $S = 3$; $\alpha = 2$ and $\beta = 3$, the pheromone matrix was initialized
 632 to $\tau_{kl}(0) = 0.5$ and the evaporation factor was set to $\rho = 0.8$. For the ACS
 633 algorithm, the parameter q_0 was initialized to 0.6 whereas for MMAS the

pheromone thresholds τ_{MIN} and τ_{MAX} are computed as suggested in [41]. Finally, the search stops once the algorithm reaches 9,000 objective function evaluations. Notice that we reduced the number of evaluations allowed for ACO-based methods, since estimating the heuristic information requires further calculations.

In the above parameter configuration, $\beta > \alpha$ since the heuristic matrix provides a suitable information source when aggregating partial rankings, which is based on the principle of greedy aggregation methods. On the other hand, the homogeneous initialization of the heuristic matrix allows guiding the search mainly based on the heuristic information at the first iterations. However, solely promoting the heuristic information may lead to stagnation or premature-convergence configurations. Aiming at preventing these undesirable states, we selected a large evaporation factor.

Moreover, we include two simpler algorithms as baselines. The first one is the *weighted Borda counting* method[13], which computes a score for each item based on its position across all observations. Next, items are arranged according to their scores. The second baseline method, baptized as the *Greedy Ant Model* (GAM), relies on a single ant in ACS that systematically only exploits the heuristic knowledge to select the next feasible state. In GAM, $\alpha = 0$, $\beta = 1$ and $q_0 = 1$. Overall, we compare our approach against two greedy methods as baseline techniques and nine evolutionary algorithms.

7.2. Numerical simulations for a real-world dataset

In this section, we evaluate the proposed methodology by using a real study case concerning to the attractiveness of companies in Belgium [50][34][27]. During the data acquisition phase, 14585 Belgian respondents (aged between 18 and 65 years old) were consulted regarding two different ranking scenarios. Both surveys were conducted by a panel of marketing experts from Randstad¹ and are summarized as follows:

- **Scenario 1.** Each respondent ranked the top-5 factors out of $M = 17$ possible factors. From this survey we obtained 14,585 partial rankings where only the top- K factors are ordered, whereas the other $M - K$ factors are tied at ranking position $K + 1$.

¹Randstad (<http://www.randstad.com>) is the second largest Human Resources (HR) provider in the world. It expanded its operations to 39 countries, representing more than 90 percent of the global HR services market.

666 • **Scenario 2.** In the second study each expert is free to select the most
667 relevant K_i factors, such that $2 \leq K_i \leq M$, therefore the number
668 of ranked elements is not necessarily homogeneous in all cases. More
669 explicitly, in the i th ranking the top- K_i factors are ordered, whereas the
670 remaining $M - K_i$ factors are ranked at the position $K_i + 1$. Moreover,
671 in this kind of partial rankings, tied factors into the top- K_i ranked
672 positions could be observed.

673 Solving these ranking aggregation problem allows characterizing the at-
674 tractiveness of companies in Belgium, that is, their ability to attract highly-
675 competent and productive employees. If workers prefer some factors (e.g.
676 comfort, salary) when they are looking for an employer, and the evaluated
677 company does not offer such features, then it is expected that more compe-
678 tent employees end up not working with that company. With this knowledge
679 at hand, the company board may improve its branding and enhance its visi-
680 bility which frequently results in better incomes. Table 3 displays the global
681 factors (ranking items) evaluated by respondents that came up after a panel
682 discussion of marketing experts.

Table 3: Global factors evaluated by each respondent during the online survey.

F_1	Financially sound
F_2	Offers quality training
F_3	Offers long-term job security
F_4	Offers international / global career
F_5	Future prospects / career opportunities
F_6	Strong management
F_7	Offers interesting jobs (job description)
F_8	Pleasant working environment
F_9	Competitive salary package
F_{10}	Good balance between life and work
F_{11}	Conveniently located
F_{12}	Strong image / pursues strong values
F_{13}	Quality products / services offered
F_{14}	Deliberately handles the environment and society
F_{15}	Uses the latest technologies / innovative
F_{16}	Provides flexible working conditions
F_{17}	Encourages diversity (age, gender, ethnicity)

683 Table 4 displays the list of expert-defined categories, which allows com-
684 puting the weight of each partial ranking as explained in Section 6. Par-
685 ticularly, we adopted the heuristic strategy for predefined categories where
686 factors are gathered according to their semantics by marketing experts.

Table 4: Categories determined by marketing experts.

	Name	Factors in the category
C_1	Salary	F_9
C_2	Stability	F_1, F_3
C_3	Future	F_2, F_5, F_4, F_7
C_4	Comfort	$F_{16}, F_{11}, F_{10}, F_8$
C_5	Status	$F_{17}, F_{14}, F_{15}, F_{13}, F_{12}, F_6$

687 Figure 1 shows the average membership degree across all categories for
688 both studies. The reader may observe that the maximal average member-
689 ship degree corresponds to the category C_2 in both scenarios. This result is
690 certainly interesting but unsurprising because Belgian people regularly have
691 well-paid jobs, and thus they are more focused on finding jobs with safer con-
692 tract terms. Therefore, the membership degrees for the “Stability” category
693 will be used to solve the weighted aggregation problem.

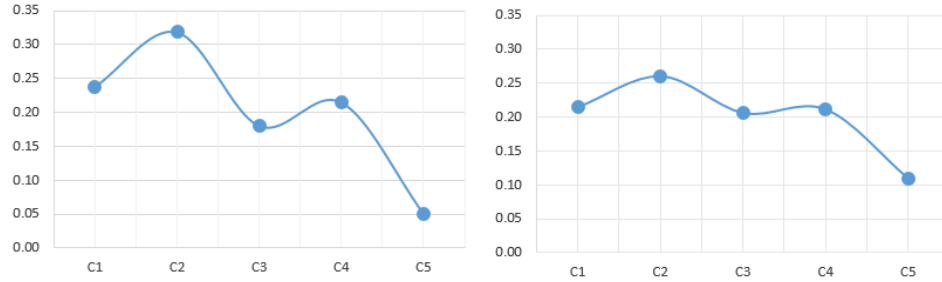


Figure 1: Average memberships degrees across predefined categories: (a) top- K aggregation scenario, and (b) top- K_i aggregation scenario.

694 In the first scenario, the differences are more evident given the lower dis-
695 persion in the selected factors (i.e., the respondents tend to include similar
696 factors into the top-5). On the other hand, in the second scenario, the disper-
697 sion increases since each respondent could select a large number of relevant
698 factors. It should be highlighted that solving the weighted Kemeny ranking

699 problem for scenarios with high dispersion is undoubtedly more complicated
700 due to the lack of consensus among respondents.

701 Figure 2 shows the performance of the selected algorithms for the top-5
702 aggregation scenario. The performance measure refers to normalized Haus-
703 dorff dissimilarity between the aggregated ranking and the partial rankings.
704 Due to the stochastic nature of evolutionary and swarm intelligence algo-
705 rithms, each record is computed from the average of 10 independent trials.
706 The reader may observe that all optimizers are capable of outperforming the
707 baseline methods, while ACS stands as the best-performing algorithm fol-
708 lowed by MMAS. Moreover, the results have shown suggest that OX2-ISM
709 is the best-performing GA-based optimizer, being ranked third overall.

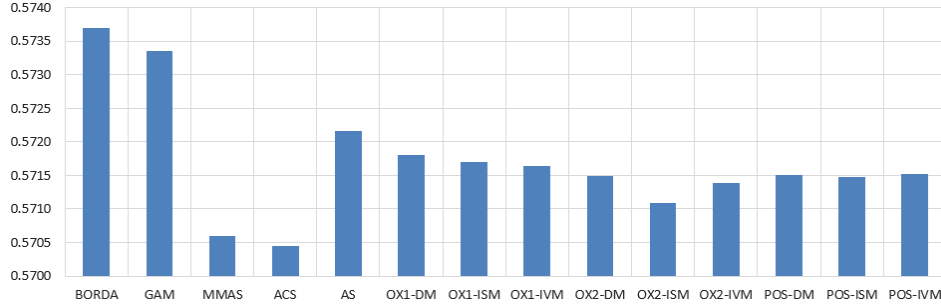


Figure 2: Performance of selected algorithms for the top-5 aggregation scenario.

710 Figure 3 reports the relative improvement rate of swarm and evolutionary
711 algorithms with regards to the two greedy methods under consideration. In
712 this aggregation scenario, the ACS and MMAS algorithms exhibit the largest
713 improvement rates.

714 Figure 4 depicts the normalized Hausdorff dissimilarity measure for the
715 top- K_i aggregation scenario. In this case, all ACO-based algorithms outper-
716 form the other approaches and ACS emerges as the top contender. The evo-
717 lutionary optimizers perform comparably among them, although algorithms
718 using the OX1 crossover operator fare slightly worst. It should be high-
719 lighted that the solutions computed by GAM are better than those produced
720 by the weighted Borda counting method. However, this greedy approach is
721 not competitive against evolutionary and swarm intelligence optimizers.

722 Figure 5 illustrates the relative improvement rate of swarm and evolution-
723 ary algorithms in comparison to the BORDA and GAM baseline methods.
724 Observe that all optimizers are capable of outperforming the baseline meth-

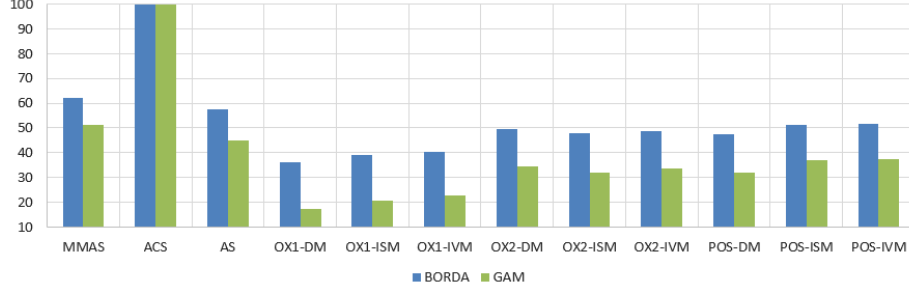


Figure 3: Improvement rate attained by the algorithms under discussion with regards to the two baseline methods for the top- K aggregation scenario.

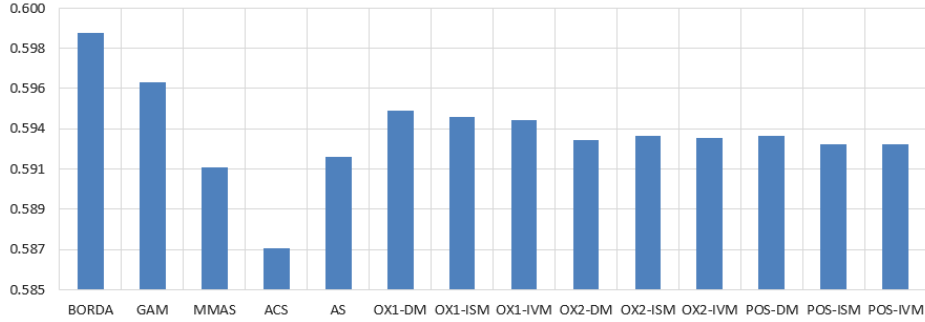


Figure 4: Performance of the selected algorithms for the top- K_i aggregation scenario.

ods and that ACS achieves the largest improvement rates.

The top- K_i datasets reveal a greater dispersion around the global (often unknown) consensus since there are fewer tied factors and finding the Kemeny ranking solution could be more challenging. Nevertheless, the inclusion of the membership degrees of each ranking to the dominant category will probably reduce the dispersion degree. This suggests that partial rankings with high membership degree to the C_2 category will comprise similar relevant factors and therefore the search problem will likely be easier to solve.

Overall, the results support the superiority of ACS and MMAS over the two greedy and nine evolutionary algorithms. The AS scheme is less competitive, which suggests that exploiting only the best solutions found by artificial ants could be convenient when aggregating partial rankings. In the next subsection, we evaluate our methodology using more generic datasets.

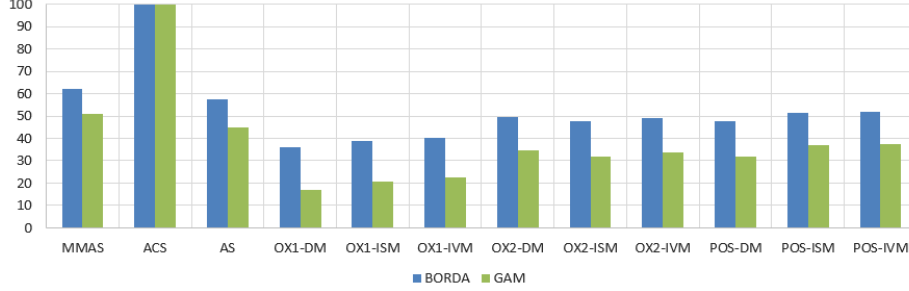


Figure 5: Improvement rate attained by the benchmarking algorithms over the two baseline methods for the top- K_i aggregation scenario.

7.3. Numerical simulations for synthetic datasets

Aiming at generalizing the above results, we adopted 190 synthetic datasets from www.PrefLib.org having different complexity in both the number of instances and factors. These datasets comprise top- K_i partial rankings where ties are allowed. The number of instances ranges from 10 to 10,335 whereas the number of factors goes from 4 to 155. The reader may observe that such datasets do not include an explicit definition of categories. Therefore, in these synthetic aggregation problems, the weight of each partial ranking is calculated according to the ratio of non-tied items as explained in section 6.

In order to verify the existence of significant differences among the suite of benchmarking algorithms, we computed the Friedman two-way analysis of variances by ranks [51]. The Friedman test is a multiple-comparisons nonparametric statistical test that detects whether at least two of the samples (in a set of $N > 2$ samples) represent populations with different median values or not. In our case, the Friedman test suggests rejecting the null hypothesis H_0 (p -value = $2.2825\text{E-}10 < 0.05$) using a confidence interval of 95%. Thus, we can conclude that there are statistically significant differences between at least two algorithms across all datasets.

Figure 6 portrays the mean ranks computed by the Friedman test. From such results, we can formalize some empirical conclusions about the performance of the methods under consideration:

- The ACS algorithm is capable of notably outperforming the remaining search methods, followed by MMAS. However, the existence of statistically significant differences between them must be verified.

- The OX1 crossover operator leads to poor performance and may not be adequate for solving the extended Kemeny ranking problem. The other evolutionary optimizers perform comparably among them, with the OX2-ISM scheme producing better results.

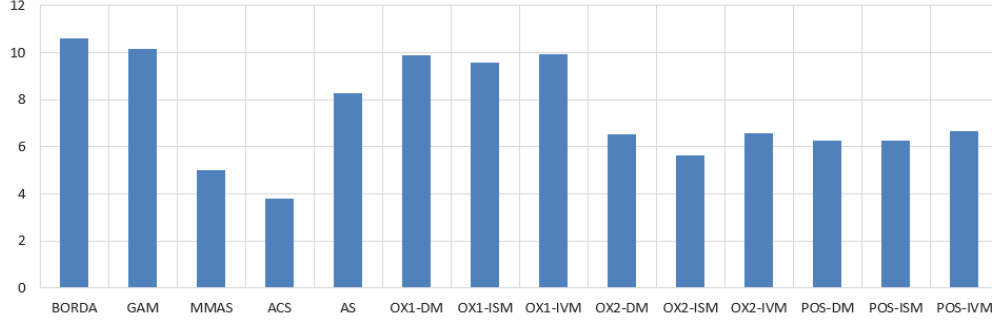


Figure 6: Mean ranks computed by the Friedman test for each algorithm across all synthetic datasets.

To further confirm the superiority of the search methods over the greedy methods, Figure 7 displays the improvement rate attained by swarm and evolutionary algorithms over these two approaches. For these synthetic datasets, ACS displays the largest improvement rates, while AS and the evolutionary methods based on the OX1 crossover operator report the worst ones.

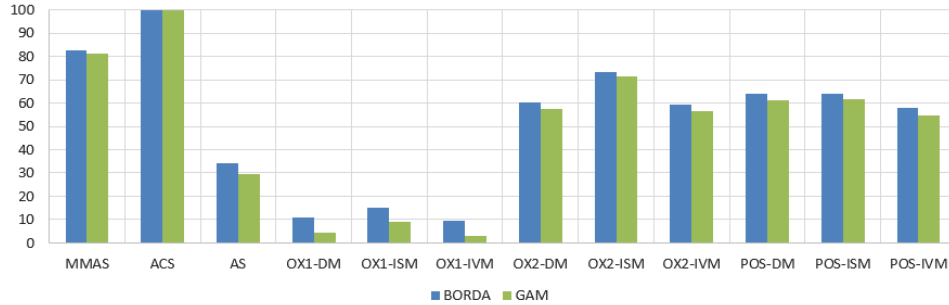


Figure 7: Improvement rate attained by the algorithms under discussion with regards to the two baseline methods for the synthetic datasets.

The last experiment is focused on determining whether the superiority of the ACS method is statistically significant or not. With this goal in mind, we use four post-hoc procedures (i.e., Bonferroni, Holm, Hochberg and Hommel)

[52] for multiple pairwise comparisons and a control method. These statistical procedures are required since in pairwise analysis, if we try to draw a conclusion involving more than one pairwise comparison, we will accumulate an error coming from its combination. Therefore, we are losing control on the Family-Wise Error Rate (FWER), defined as the probability of making one or more false discoveries among all the hypotheses when performing multiple pairwise tests [52]. Table 5 reports the unadjusted p -value and the adjusted p -values associated with each pairwise comparison using the best-performing algorithm (ACS) as the control method.

Table 5: Post-hoc procedures for pairwise comparisons using ACS as the control algorithm.

Algorithm	p -value	Bonferroni	Holm	Hochberg	Hommel
BORDA	1.8939E-56	2.4621E-55	2.4621E-55	2.4621E-55	2.4621E-55
GAM	2.0719E-49	2.6935E-48	2.4863E-48	2.4863E-48	2.4863E-48
OX1-IVM	1.8769E-46	2.4399E-45	2.0645E-45	2.0645E-45	2.0645E-45
OX1-DM	3.0934E-45	4.0214E-44	3.0934E-44	3.0934E-44	3.0934E-44
OX1-ISM	4.9068E-41	6.3789E-40	4.4161E-40	4.4161E-40	4.4161E-40
AS	2.0687E-25	2.6893E-24	1.6550E-24	1.6550E-24	1.6550E-24
POS-IVM	2.5414E-11	3.3038E-10	1.7789E-10	1.7789E-10	1.7789E-10
OX2-IVM	1.2108E-10	1.5741E-9	7.2652E-10	7.2652E-10	7.2652E-10
OX2-DM	2.8047E-10	3.6461E-9	1.4023E-9	1.4023E-9	1.4023E-9
POS-DM	1.1409E-8	1.4831E-7	4.5636E-8	3.9519E-8	3.4227E-8
POS-ISM	1.3173E-8	1.7125E-7	4.5636E-8	3.9519E-8	3.9519E-8
OX2-ISM	2.2062E-5	2.8681E-4	4.4125E-5	4.4125E-5	4.4125E-5
MMAS	0.0055	0.0725	0.0055	0.0055	0.0055

All post-hoc procedures reject the null hypothesis for a 5% significance level (corresponding to the 95% confidence interval). Only the Bonferroni procedure accepts the conservative hypothesis for the ACS-MMAS pair. This allows claiming the statistical superiority of the ACS optimizer when tested with the synthetic datasets used for simulations. From these results we can state that the strong elitism embedded into ACS/MMAS is a determinant aspect when aggregating weighted partial rankings.

8. Concluding remarks

In this paper we addressed the weighted aggregation of top- K and top- K_i partial rankings by extending the Kemeny ranking problem. To do that,

we proposed a search method rooted on Ant Colony Optimization that automatically estimates the heuristic information from the rank population. Moreover, we discussed two strategies to compute the weight of each partial ranking. In the first case, the weight is computed from the fuzzy membership degree of the target instance to a set of predefined categories (i.e., information granules resulting from a partition of the whole item set). If such categories are not available, then the weight is calculated as the ratio of non-tied items. This heuristic assumes that partial rankings having fewer number of tied factors are more informative when aggregating partial rankings. The reader may observe that other alternatives to determine the weights could be explored.

During the simulations, we compared the performance of our approach against two (greedy) baseline methods and nine genetic-algorithm-based implementations in presence of both a real-world problem and 190 synthetic datasets. The results have shown that the ACS algorithm is capable of significantly outperforming the remaining techniques, followed by the MMAS algorithm. This finding could be ascribed to the elitist approach of ACS and MMAS in conjunction with the proposed strategy for estimating the heuristic information. Moreover, we observed that the OX1 crossover operator reports poor performance; therefore, it may not be suitable to solve the Kemeny ranking problem. As a future work, we will focus on extending the proposed approach to more generic aggregation scenarios.

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