

## The Distribution of N-Grams

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# THE DISTRIBUTION OF N-GRAMS

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## ABSTRACT

N-grams are generalized words consisting of N consecutive symbols, as they are used in a text. This paper determines the rank-frequency distribution for redundant N-grams. For entire texts this is known to be Zipf's law (i.e. an inverse power law). For N-grams, however, we show that the rank (r)-frequency distribution is

$$P_N(r) = \frac{C}{(\psi_N(r))^\beta},$$

where  $\psi_N$  is the inverse function of  $f_N(x) = x \ln^{N-1} x$ . Here we assume that the rank-frequency distribution of the symbols follows Zipf's law with exponent  $\beta$ .

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## I. Introduction

N-grams are generalized words consisting of  $N$  consecutive symbols. N-grams are a very important study topic as they have applications in indexing, information retrieval, error correction, text compression, language identification, subject classification and even speech recognition.

Consider a text (in principle any kind of text) and fix  $N \in \mathbb{N}$ . Usually  $N=3, 4, 5$  or  $6$  but this is not important here. In this text, using a “window” technique, one determines a (long) list of occurring N-grams (see e.g. Robertson and Willett (1998)). Each type of N-gram is a coordinate in a vector, representing the text under study and the number appearing in this coordinate is the number of times this N-gram appears in the text. In this way texts can be compared by using similarity measures between the vectors of any two of these texts. In this way texts can be clustered when they are similar and this information is important in indexing. Classical similarity measures are : the Dice index (Robertson and Willett (1998)), the Jaccard index (Hamers et al. (1989), Rousseau (1998)) or Salton’s cosine formula (Salton and Mc Gill (1987)) but other similarity measures can be used.

This procedure can also be used on key words, hence associating with them similar key words (by applying a threshold on the similarity measure) to be used in indexing. Indexing and information retrieval (IR) are dual concepts (Egghe and Rousseau (1997)) and hence N-grams have applications in IR : clustered documents (as described above) are also retrieved and queries can be expanded by determining (as described above) similar key words. This technique is a considerable extension of the classical truncation technique.

Error correction can be done if a certain word is not recognized in a system (e.g. a word does not belong to the dictionary) and hence the word is considered to be in error. If not it is added as such ; in the other case similar words (determined by N-grams-similarity) are proposed for correction.

Language identification and subject classification are also cluster-based techniques, based on the fact that languages or topics (in a language) are characterized by their N-grams distributions. Similarities can be calculated and clusters (or other multivariable statistical methods) can be determined, hence presenting “atlases” of languages or of topics (or scientific disciplines). In the same way can N-grams be used in the understanding of natural language.

Text compression is obtained by attaching heavily used N-grams to unused bytes which are reserved for single symbols (since the set of single symbols (hence 1-grams) usually is not an integer power of 2).

Finally, N-grams of phonemes or of words (not of letters) can be used in speech recognition. We refer the reader to the review articles Cohen (1995), Damashek (1995), Robertson and Willett (1998) and Yannakoudakis, Tsomokos and Hutton (1990) and to the many references therein. Note that all these techniques are language or subject independent (contrary to more classical ones, e.g. truncation, indexing, ...). It is however noted in the above references that these techniques perform better in Asian languages, because of their specific structure (symbols are entire syllables and hence truncation - to mention only this example - works not so well in these languages).

So, this should convince the reader that the study of N-grams is very important. One element in such a study is their rank-frequency distribution, i.e. fixing N, determine a relationship between the number of times a certain N-gram appears and its rank  $r$  (ranked in decreasing order of appearance). Hence we are interested in the N-gram analogue of the law of Zipf which is known to be valid in ordinary texts (which can be considered as texts consisting of N-grams but with variable N, i.e. the words of different lengths). As is the case with the law of Zipf, such a distribution determines the compressive power of N-grams. Indeed, the more concentrated (see Egghe and Rousseau (1990a,b, 1991) N-grams are (i.e. the more unequal their distribution is) the more effective compression that can be obtained.

Zipf's law is well-known. It states that

$$x = P(r) = \frac{C}{r^\beta}, \quad (1)$$

where  $r$  is the rank of a word (decreasing order of occurrence),  $P(r)$  is the fraction describing how often this word occurs and  $\beta$  and  $C$  are constants. Otherwise stated

$$r = \frac{C^{\frac{1}{\beta}}}{x^{\frac{1}{\beta}}}, \quad (2)$$

also a (decreasing) power law.

Is such a law also valid if one only considers  $N$ -grams ( $N$  fixed)? In the next section we prove that for redundant  $N$ -grams (see the next section for definitions)

$$r = D \frac{C^{\frac{N}{\beta}}}{x^{\frac{1}{\beta}}} \ln^{N-1} \left( \frac{C^{\frac{N}{\beta}}}{x^{\frac{1}{\beta}}} \right) \quad (3)$$

hence the composition of a “classical” Zipf law and the function

$$f_N(y) = y \ln^{N-1} y. \quad (4)$$

In other words, the exact law of Zipf is not valid in the case of redundant  $N$ -grams ( $N$  fixed). In the proof we assume that the symbols’ rank-frequency distribution is Zipfian, a fact that can be accepted in a first approximation (also this is discussed in the next section).

This result is similar to a result obtained in Egghe (1999) on the distribution of  $N$ -word phrases but the present proof is a simplification of the argument given there and, in addition, is more general so that it also comprises  $N$ -grams distributions.

We note that, in our arguments for fixed  $N$ , we need less assumptions as e.g. in the arguments of Li (1992) or of Mandelbrot, explaining his law for words in texts, based on symbols (see Mandelbrot (1977) or Egghe and Rousseau (1990a) where the proof of Mandelbrot is more accessible). There one assumes that all symbols have an equal chance to occur. This is

certainly not the case (a trivial fact) e.g. illustrated by Dewey's table appearing in Heaps (1978), p. 200). Using that all symbols have an equal chance to occur also implies that, the longer a word is the lesser it is used! This is certainly not the case : words as the, an, it, ... are certainly used more heavily than letters such as j, z or q! Our model uses the real probabilities of the symbols and this distribution has a parameter influence on the final result. In Perline (1996) one also presents an argument using real symbol probabilities (although completely different from ours) but this argument seems to be in error (Troll and beim Graben (1998)).

## II. The rank-frequency distribution of N-grams

In his analysis of theoretical texts, in which words are formed by symbols and blanks, Mandelbrot (1977) assumes the following simplifications of real life :

1. Symbols (mostly letters) follow independently in words,
2. All words (i.e. N-grams,  $N \in \mathbb{N}$ ) are used,
3. All symbols have an equal probability of occurrence.

None of these assumptions is correct and this is so in any text and in any language. Also, condition 3 has the consequence that, the longer a word is, the less it is used. This is false in any language. Take the example of English : words like "the", "an", "in", "of" and so on are used more than the 1-letter words "x" or "q" or "z". Also to look at this article, it is clear that the word "distribution" is used more than the word "cat".

So the novelty of the argument given in this article is - first of all - allowing the symbols to have a realistic distribution of occurrence, i.e. rank-frequency distribution. We have inspected the following sources :

- (i) The data of Dewey (as given in Heaps (1978), p. 200) on the occurrence of the 26 letters in English. The graph is given in Fig. 1.

- (ii) The data, as extracted from one of my reports of the LUC-library management committee (in Flemish). The graph is given in Fig. 2.
- (iii) The data extracted from the Website of Beckman (see Beckman (1999)) on the occurrence of more than 13,000 Chinese symbols. The graph is given in Fig. 3a and b for the first 300 and first 100 characters, respectively.
- (iv) The 3 texts analysed by Yannakoudakis et al. (1990). One text is a book in modern Greek (here we have 24 letters) ; the two others are text books in English (one on the history of the Mormon church and one on database management systems). The graph of the Greek text is given in Fig. 4. The others are very similar.

It is clear that, replacing the uniform distribution (as is the case in simplification (3)) by a decreasing power law of the type

$$P_1(r) = \frac{C}{r^\beta} \quad (5)$$

is a good first order assumption.

Note that (5) is equivalent with

$$r = \frac{C^{\frac{1}{\beta}}}{P_1(r)^{\frac{1}{\beta}}} \quad (6)$$

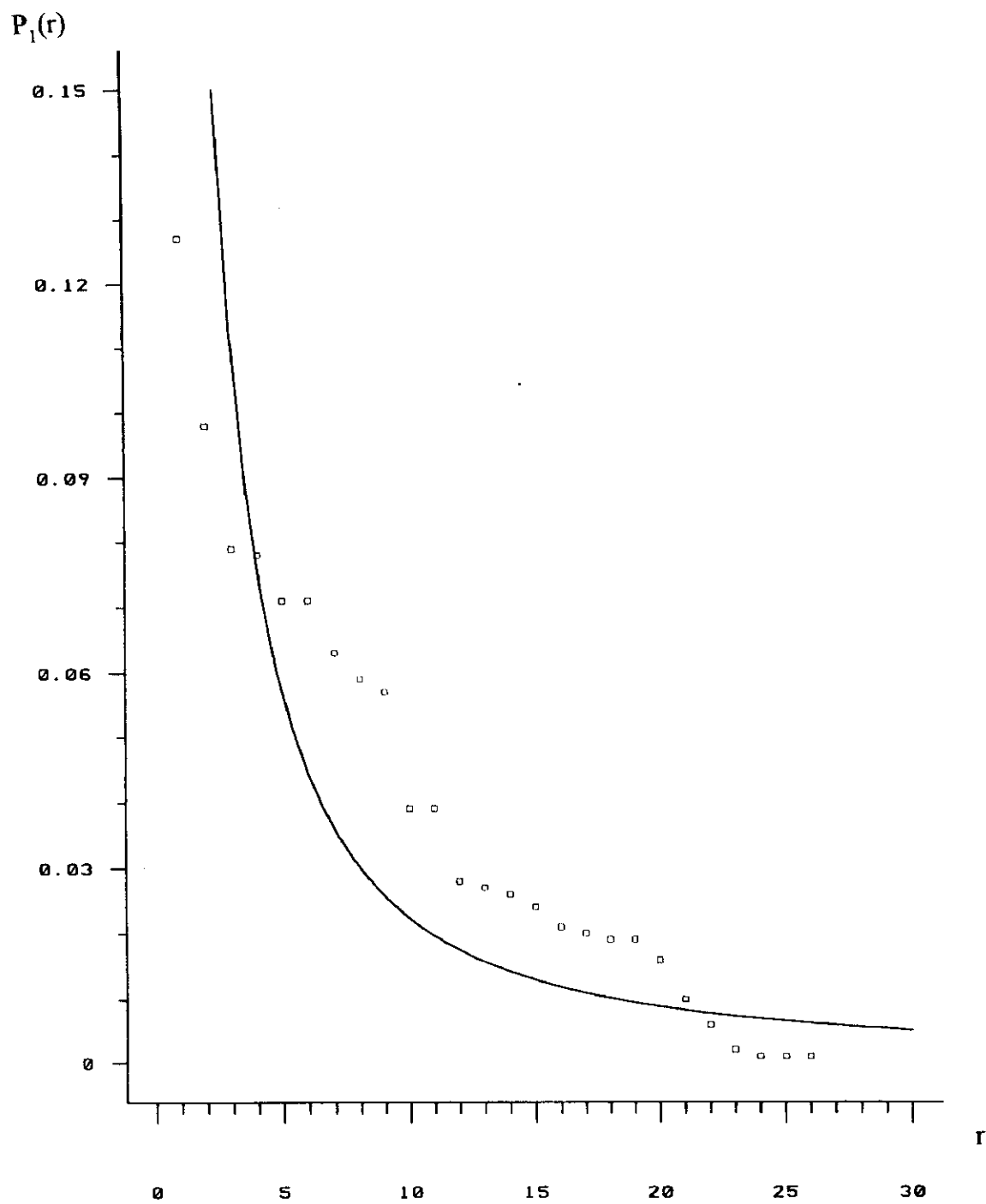


Fig. 1 Letter occurrence, Dewey data (English)



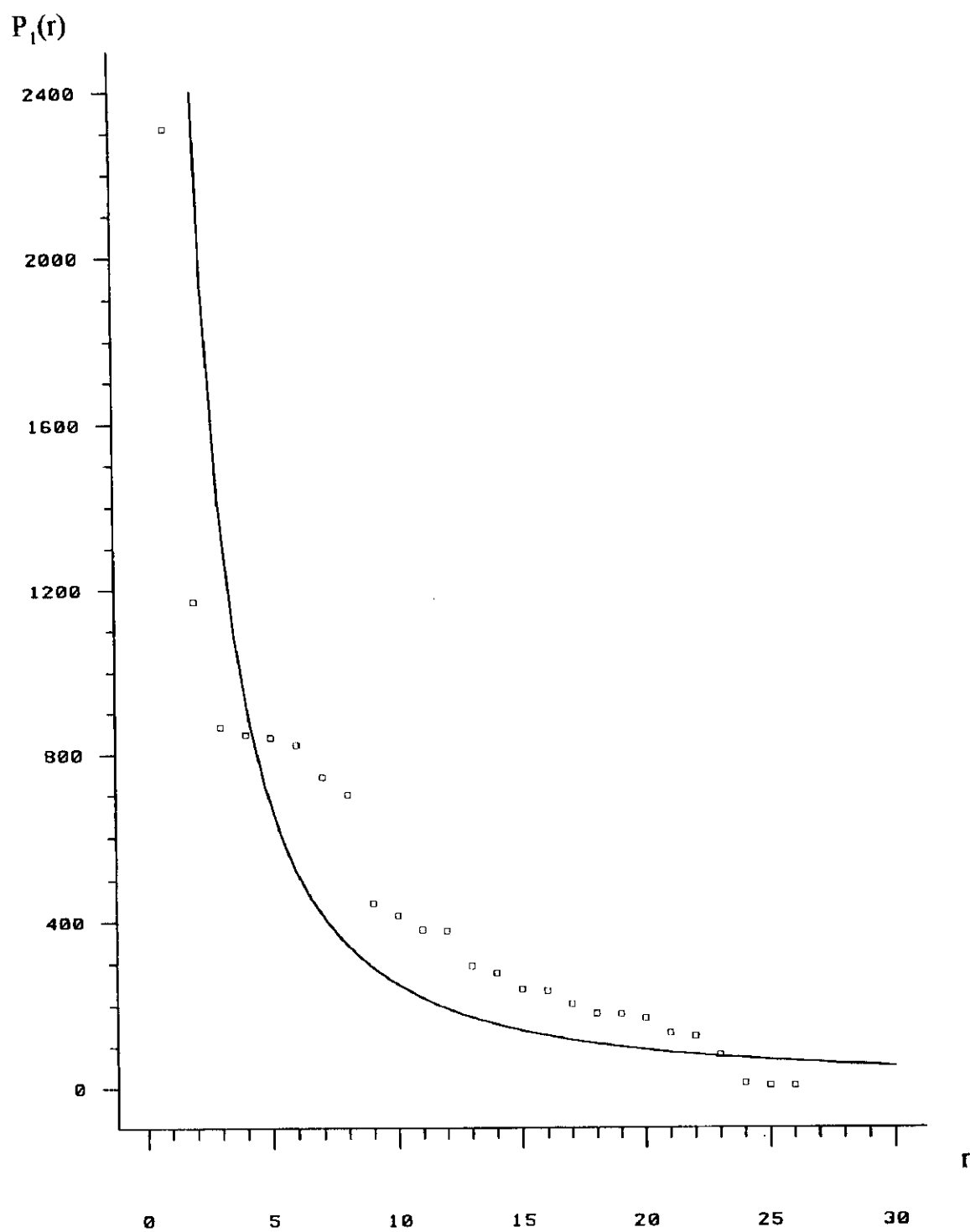


Fig. 2

Letter occurrence, LUC data (Flemish)

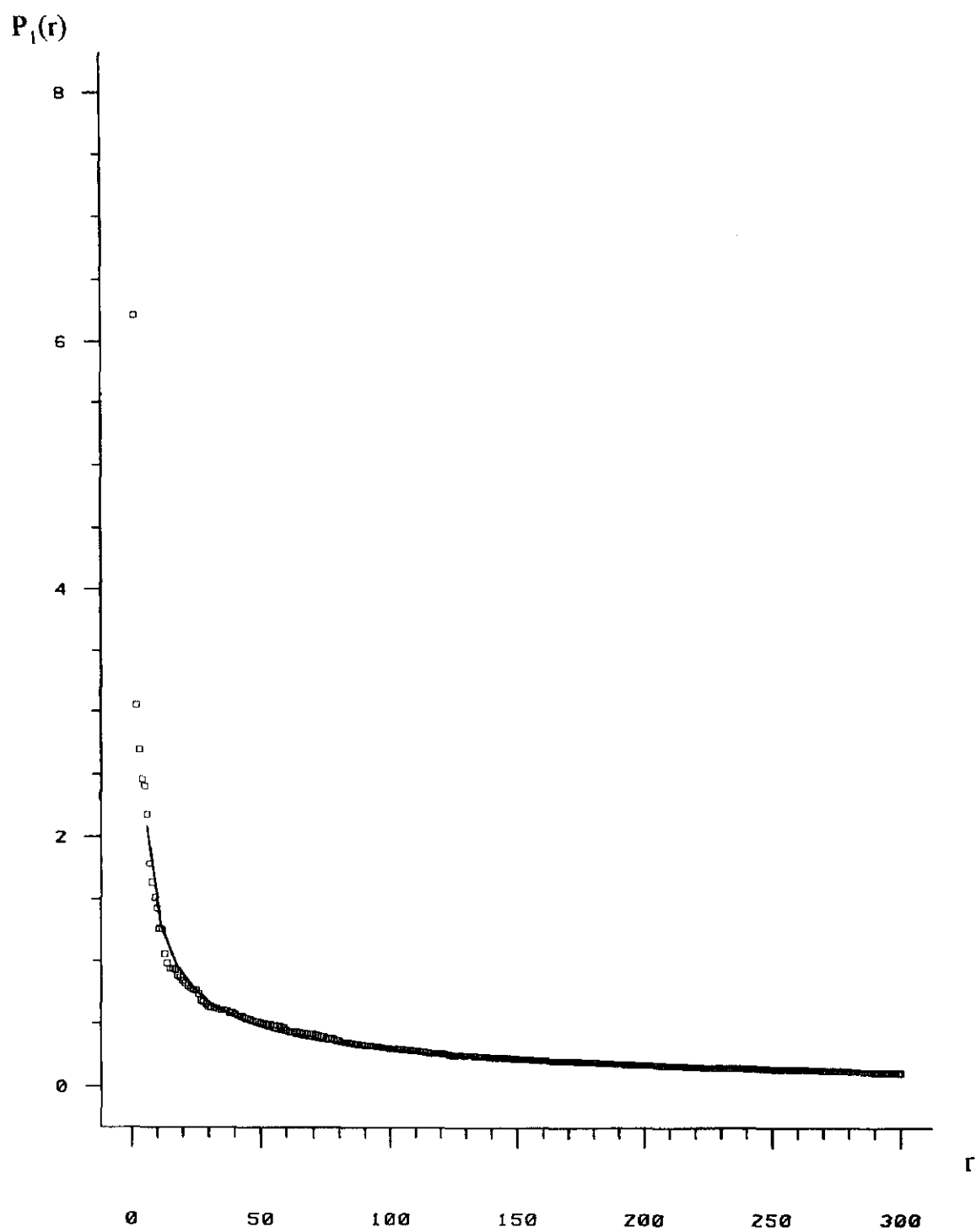


Fig. 3a Occurrence of the 300 mostly used Chinese characters

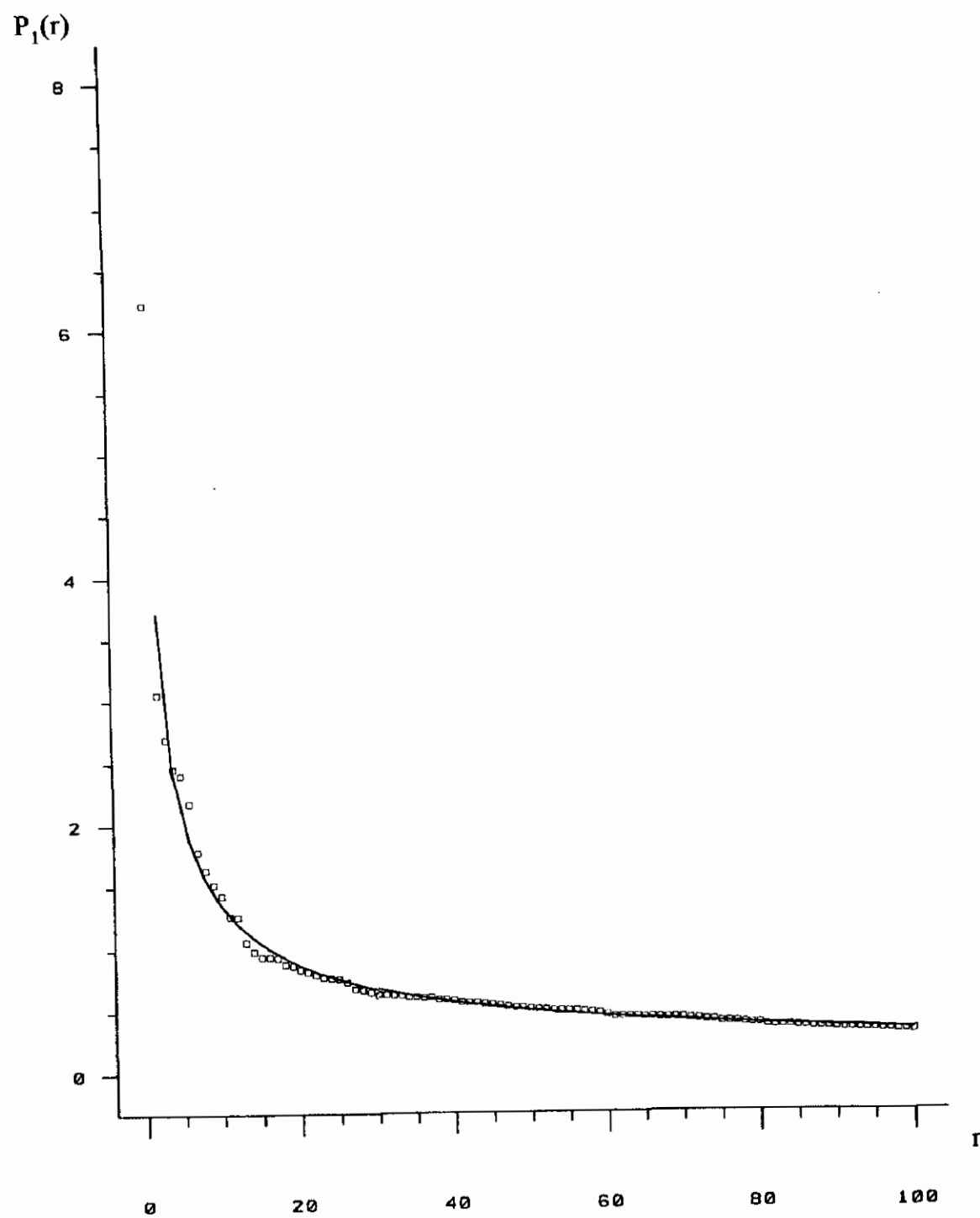


Fig. 3b Occurrence of the 100 mostly used Chinese characters

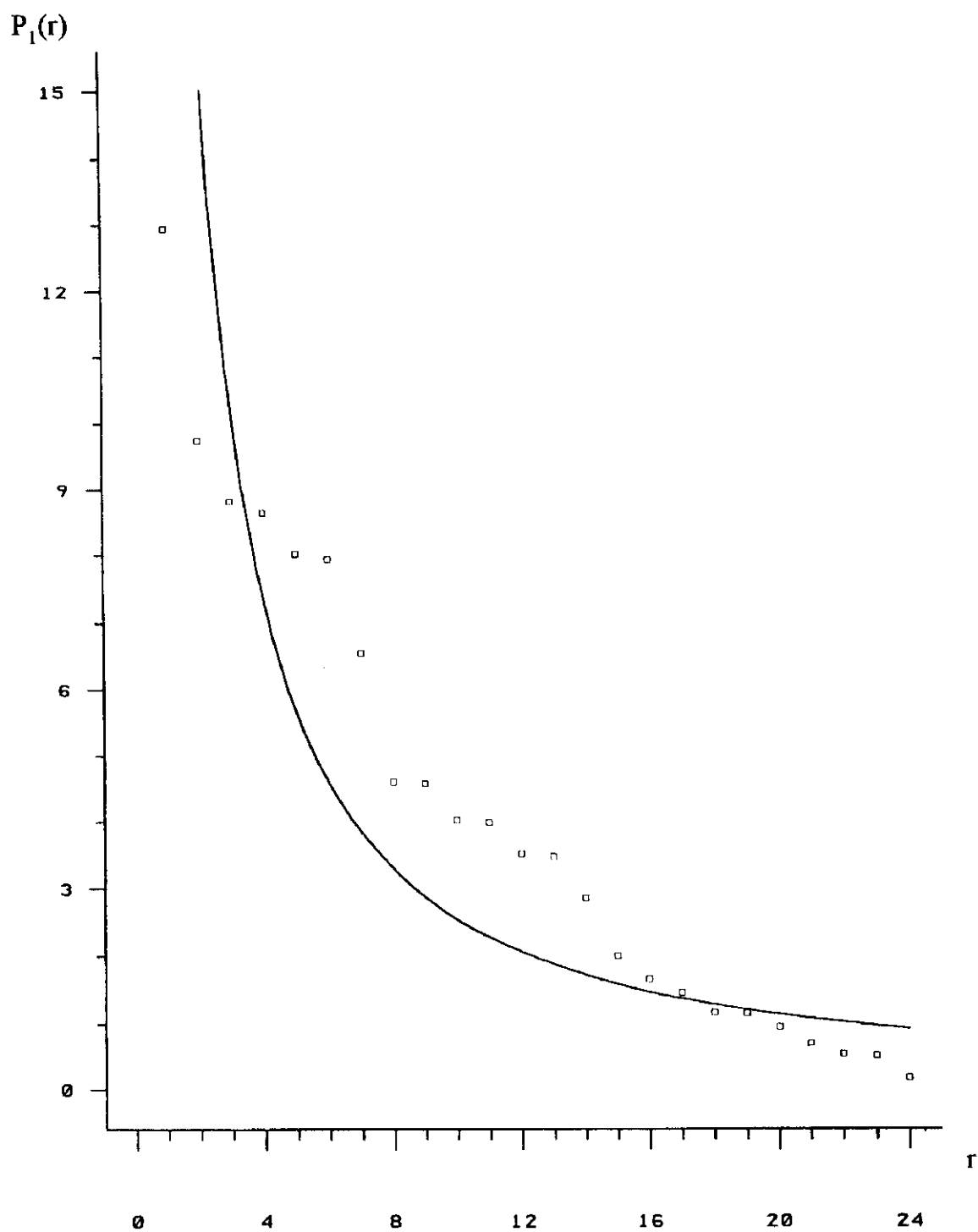


Fig. 4 Letter occurrence, ARG01 data (Greek)

It is even almost perfect for the Chinese example (and we assume that this is also the case for the other Asian languages as well). Note that we remarked in the introduction that N-gram techniques work best in Asian languages.

This assumption also solves the simplification (2). Indeed, by allowing different probabilities for symbols in texts, we automatically add an extra diversification in the occurrence of words, exactly as they should appear. As to simplification (1) we note that, in general texts, it is certainly wrong but in the case studied here, it is true. This can be seen as follows. We will use N-grams generated by a window in an ordinary text. We will use the redundant way of generation (see e.g. Willett (1979)), the mostly used way to generate N-grams (see Robertson and Willett (1998)). To state the exact meaning of redundant generation of N-grams we show the generation from the word SYMBOL by 2- and 3-grams

redundant

2-Gram

\*S SY YM MB BO OL L\*

3-Gram

\*\*S \*SY SYM YMB MBO BOL OL\* L\*\*

non-redundant

2-Gram

SY MB OL

3-Gram

SYM BOL

(in the non-redundant case, incomplete N-grams are not used).

As is obvious from the redundant generation of N-grams, EVERY symbol (appearing in a word) appears once in the  $i^{\text{th}}$  place of an N-gram ( $i=1,\dots,N$ ). Hence, the occurrence of the symbols in N-grams is independent of their place and of the symbols before them. In other words, the symbol distribution in any place  $i$  ( $i=1,\dots,N$ ) in an N-gram is the same as the overall symbol distribution in the text. In mathematical notation : let  $r_i$  be the rank of the symbol in the  $i^{\text{th}}$  place in the N-gram ( $i=1,\dots,N$ ). Then, because of the above, the conditional probability to have  $r_i$  at the  $i^{\text{th}}$  place, given  $r_j$  at the  $j^{\text{th}}$  place ( $j=1,\dots,i-1$ ) is

$$\begin{aligned} & P(r_i|r_1, \dots, r_{i-1}) \\ &= P_1(r_i), \end{aligned} \tag{7}$$

where  $P_1(\cdot)$  denotes the symbol distribution in the text. For this distribution, we assume (5) (as explained above). The symbol set is assumed to be large enough so that also other symbols than letters can be used, or, in Asian languages, also all symbols can be considered). Since the alphabets in the latter cases are so high we can easily assume it to be infinite. In formula (5) we used the symbol  $P_1$  denoting probabilities for symbols, i.e. 1-grams. Our task is,  $N \in \mathbb{N}$  being fixed, to determine the distribution  $P_N(r)$ , denoting the probability for the  $N$ -gram on rank  $r$  to occur in the text. As ever, ranks are given according to the decreasing order of occurrence.

We have the following result.

**Theorem II.1.** Let  $N \in \mathbb{N}$  be fixed and assume (5) to be valid. Let  $P_N(r)$  denote the probability (fraction) of occurrence of the  $N$ -gram on rank  $r$ . Then

$$r = \frac{C^{\frac{N}{\beta}}}{P_N(r)^{\frac{1}{\beta}}} \frac{\ln^{N-1} \left( C^{\frac{N}{\beta}} / P_N(r)^{\frac{1}{\beta}} \right)}{(N-1)!} \tag{8}$$

, hence

$$r = \frac{1}{(N-1)!} f_N(y), \tag{9}$$

where  $f_N(y) = y \ln^{N-1}(y)$  and  $y = \frac{C^{\frac{N}{\beta}}}{P_N(r)^{\frac{1}{\beta}}}$

(Note that for  $N=1$ , we refind (6)). Otherwise stated, there exist constants  $\alpha_N$  and  $D_N$  such that

$$P_N(r) = \frac{D_N}{(\Psi_N(\alpha_N r))^{\beta}}, \tag{10}$$

where  $\psi_N = f_N^{-1}$ . For large  $r$  this reduces to : there exists a constant  $E_N$  such that

$$P_N(r) \approx \frac{E_N}{(\psi_N(r))^\beta}, \quad (11)$$

hence a power law, but not of  $r$  but of  $\psi_N(r)$ . Note that the same exponent  $\beta$  (as in (5)) appears.

The proof is rather long, based on (5) and (7) and given in the Appendix. Nevertheless it is a simplification of the proof given in Egghe (1999) on the distribution of N-word phrases. The proof there can equally be simplified in this way. Note however that the result obtained in Egghe (1999) is an approximation of reality by assuming that the words in an N-word phrase occur independently from each other. The result obtained here for N-grams is exact since the letters indeed occur independently from each other (by (7)).

### III. Conclusions and open problems.

We have shown that, under more realistic assumptions than in Mandelbrot (1977), where we assume that the symbols follow a power law (instead of a uniform distribution) and where they occur independently of each other, that the N-gram distribution is a power law of  $\psi_N(r)$ , where  $\psi_N$  is the inverse of the function

$$f_N(x) = x \ln^{N-1} x.$$

This result is new and is obtained without any simplification (as mentioned in section II) or approximation.

As open problem we can formulate :

Establish a theory in which general texts (words with arbitrary lengths) can be treated.

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## Appendix

### Proof of Theorem II.1

It is clear from (7) that, for every  $x \in ]0, 1[$ , if we put

$$r = \#\{(r_1, \dots, r_N) \mid P_1(r_1) \dots P_1(r_N) \geq x\} \quad (A1)$$

we have that  $x = P_N(r)$ . Indeed : the probability for an N-gram for which the  $i^{\text{th}}$  symbol has rank  $r_i$  ( $i=1, \dots, N$ ) is  $P(r_1 | r_2, \dots, r_N) P(r_2 | r_1, \dots, r_N) \dots P(r_N | r_1, \dots, r_{N-1}) P(r_1) = P(r_1) P(r_2 | r_1) \dots P(r_N | r_1, \dots, r_{N-1}) P(r_1)$ , by (7).

Note that in this paper  $N$  is fixed. We will consider  $x \in ]0, 1[$  as a continuous variable, hence using continuous distributions  $P_N(r)$ ,  $r \geq 1$ . In this sense (A1) should be interpreted as the volume of the N-dimensional set

$$\{(r_1, \dots, r_N) \mid P_1(r_1) \dots P_1(r_N) \geq x\}.$$

The inequality  $P_1(r_1) \dots P_1(r_N) \geq x$  leads to, using (5)

$$\frac{C^N}{(r_1 \dots r_N)^\beta} \geq x \quad (A2)$$

hence

$$r_1 \dots r_N \leq \frac{C^{\frac{N}{\beta}}}{x^{\frac{1}{\beta}}} \quad (A3)$$

To simplify this temporary situation denote

$$a = \frac{C^{\frac{N}{\beta}}}{x^{\frac{1}{\beta}}} \quad (A4)$$

(A3) implies

$$1 \leq r_1 \leq \frac{a}{r_2 \cdots r_N} . \quad (\text{A5})$$

This will give us the number of possible  $r_1$ s, but dependent on the number of different  $r_2, \dots, r_N$ s that are possible. This will be determined now. (A5) implies

$$1 \leq r_2 \leq \frac{a}{r_3 \cdots r_N} . \quad (\text{A6})$$

(A6) implies

$$1 \leq r_3 \leq \frac{a}{r_4 \cdots r_N} \quad (\text{A7})$$

and so on, until

$$1 \leq r_{N-1} \leq \frac{a}{r_N} \quad (\text{A8})$$

and

$$r_N \leq a . \quad (\text{A9})$$

We have by (A1), (A5)-(A9) :

$$r = a \int_{r_N=1}^{r_N=a} \frac{dr_N}{r_N} \int_{r_{N-1}=1}^{r_{N-1}=\frac{a}{r_N}} \frac{dr_{N-1}}{r_{N-1}} \cdots \int_{r_2=1}^{r_2=\frac{a}{r_3 \cdots r_N}} \frac{dr_2}{r_2} \quad (\text{A10})$$

After a long, inductive but easy calculation, (A10) resolves in :

$$r = \frac{a \ln^{N-1} a}{(N-1)!} \quad (\text{A11})$$

By (A4) and the fact that  $x = P_N(r)$ , formula (7) follows. Hence also (8). (A11) now implies

$$\frac{C^{\frac{N}{\beta}}}{x^{\frac{1}{\beta}}} \ln^{N-1} \left( \frac{C^{\frac{N}{\beta}}}{x^{\frac{1}{\beta}}} \right) = (N-1)! r.$$

Hence

$$\frac{C^{\frac{N}{\beta}}}{x^{\frac{1}{\beta}}} = \psi_N((N-1)! r), \quad (\text{A12})$$

where  $\psi_N = f_N^{-1}$  and where  $f_N(y) = y \ln^{N-1}(y)$ . So

$$x = P_N(r) = \frac{D_N}{(\psi_N(\alpha_N r))^{\beta}}, \quad (\text{A13})$$

for  $D_N = C^N$  and  $\alpha_N = (N-1)!$ , yielding (9). For large  $r$ , we proceed as follows from (A11) :

$$\begin{aligned} r^{\frac{1}{N-1}} &= \left( \frac{a}{(N-1)!} \right)^{\frac{1}{N-1}} \ln a \\ &= \left( \frac{a}{(N-1)!} \right)^{\frac{1}{N-1}} \ln \left( \frac{a}{(N-1)!} \right) - \left( \frac{a}{(N-1)!} \right)^{\frac{1}{N-1}} \ln \left( \frac{1}{(N-1)!} \right) \end{aligned}$$

But, for  $r$  high (equivalently  $x$  small),  $a$  is high by (A4). Hence

$$r^{\frac{1}{N-1}} \approx \left( \frac{a}{(N-1)!} \right)^{\frac{1}{N-1}} \ln \left( \frac{a}{(N-1)!} \right).$$

So

$$r \approx \frac{a}{(N-1)!} \ln^{N-1} \left( \frac{a}{(N-1)!} \right).$$

Hence

$$\frac{a}{(N-1)!} \approx \psi_N(r),$$

by definition of  $r$ . Consequently (10) follows for

$$E_N = \frac{C^N}{((N-1)!)^\beta},$$

using again that  $a = \frac{C^{\frac{N}{\beta}}}{x^{\frac{1}{\beta}}}$ .  $\square$