

Structural Characterizations of the Semantics of Xpath as Navigation  
Tool on a Document

Peer-reviewed author version

GYSENS, Marc; Paredaens, Jan; Van Gucht, Dirk & Fletcher, George H.L. (2006)  
Structural Characterizations of the Semantics of Xpath as Navigation Tool on a  
Document. In: Vansummeren, Stijn (Ed.) Proceedings of the 25th ACM  
SIGMOD-SIGACT-SIGART Symposium on Principles of Database Systems (PODS  
2006). p. 318-327..

Handle: <http://hdl.handle.net/1942/1411>

# Structural Characterizations of the Semantics of XPath as Navigation Tool on a Document

Marc Gyssens  
Hasselt University &  
Transnational University of  
Limburg  
marc.gyssens@uhasselt.be

Jan Paredaens  
University of Antwerp  
jan.paredaens@ua.ac.be

Dirk Van Gucht  
George H.L. Fletcher  
Indiana University,  
Bloomington  
vgucht@cs.indiana.edu  
gefletch@cs.indiana.edu

## ABSTRACT

Given a document  $D$  in the form of an unordered labeled tree, we study the expressibility on  $D$  of various fragments of XPath, the core navigational language on XML documents. We give characterizations, in terms of the structure of  $D$ , for when a binary relation on its nodes is definable by an XPath expression in these fragments. Since each pair of nodes in such a relation represents a unique path in  $D$ , our results therefore capture the sets of paths in  $D$  definable in XPath. We refer to this perspective on the semantics of XPath as the “global view.” In contrast with this global view, there is also a “local view” where one is interested in the nodes to which one can navigate starting from a particular node in the document. In this view, we characterize when a set of nodes in  $D$  can be defined as the result of applying an XPath expression to a given node of  $D$ . All these definability results, both in the global and the local view, are obtained by using a robust two-step methodology, which consists of first characterizing when two nodes cannot be distinguished by an expression in the respective fragments of XPath, and then bootstrapping these characterizations to the desired results.

## Categories and Subject Descriptors

H.2.3 [Database Management]: Languages—*query languages*

## General Terms

Languages, Theory

## Keywords

XPath, expressibility, definability

## 1. INTRODUCTION

XPath is a simple language for navigation in XML documents which is at the heart of standard XML transformation languages and other XML technologies [4].

XPath can be viewed as a query language in which an expression associates to every document a binary relation on its nodes

representing all navigation paths in the document defined by that expression [3, 11, 18]. From that query-level perspective, several natural semantic issues have been investigated in recent years for various fragments of XPath. These include expressibility, closure properties, and complexity of evaluation [3, 12, 18], as well as decision problems such as satisfiability, containment, and equivalence [2, 19].

Alternatively, we can view XPath as a navigational tool on a particular given document, and study expressiveness issues from this document-level perspective. (A similar duality exists in the relational database model, where Bancilhon [1] and Paredaens [21] considered and characterized expressiveness at the instance level, which, subsequently, Chandra and Harel [7] contrasted with expressiveness at the query level.)

In this setting, our goal is to characterize, for various natural fragments of XPath, when a binary relation on the nodes of a given document (i.e., a set of navigation paths) is definable by an expression in the fragment.

To achieve this goal, we develop a robust two-step methodology. The first step consists of characterizing when two nodes in a document cannot be distinguished by an expression in the fragment under consideration. It turns out for those fragments we consider that this notion of expression equivalence of nodes is equivalent to an appropriate generalization of bisimilarity. The second step of our methodology then consists of bootstrapping this result to a characterization for when a binary relation on the nodes of a given document is definable by an expression in the fragment (in the sense of the previous paragraph).

We refer to this perspective on the semantics of XPath at the document level as the “global view.” In contrast with this global view, there is also a “local view” which we consider. In this view, one is only interested in the nodes to which one can navigate starting from a particular given node in the document under consideration. From this perspective, a set of nodes of that document can be seen as the end points of a set of paths starting at the given node. For each of the XPath fragments considered, we characterize when such a set represents the set of *all* paths starting at the given node defined by some expression in the fragment. These characterizations are derived from the corresponding characterizations in the “global view,” and turn out to be particularly elegant in the important special case where the starting node is the root.

In this paper, we study four XPath fragments. The most expressive among them is the *XPath-algebra* which permits the self, parent, and child operators, predicates, compositions, and the boolean operators union, intersection, and difference. (Since we work at the document level, i.e., the document is given, there is no need to

Permission to make digital or hard copies of all or part of this work for personal or classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation on the first page. To copy otherwise, to republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee.

PODS'06, June 26–28, 2006, Chicago, Illinois, USA.  
Copyright 2006 ACM 1-59593-318-2/06/0003 ...\$5.00.

include the ancestor and descendant operators as primitives.) We also consider the *core XPath-algebra*, which is the XPath-algebra without intersection and difference at the expression level. The core XPath-algebra is the adaptation to our setting of Core XPath of Gottlob et al. [11]. Of both of these algebras, we consider the fragments without the parent operator, called the *downward XPath-algebra* and *downward core XPath-algebra*, respectively.

The robustness of the characterizations provided in this paper is further strengthened by their feasibility. As discussed in Section 8, the global and local definability problems for each of the XPath fragments are decidable in polynomial time. This feasibility hints towards efficient partitioning and reduction techniques on both the set of nodes and the set of paths in a document. Such techniques may be fruitfully applied towards document compression [6], access control [9], and designing indexes for query processing [10, 14, 20, 22].

The remainder of this paper is organized as follows. In Section 2, we formally define the four XPath fragments as well as expression equivalence of nodes, and introduce some terminology. In Section 3, we propose our two-step methodology by applying it to both downward fragments of XPath, because these allow the simplest exposition. In particular, it will turn out that both fragments are equivalent, and that, in these cases, expression equivalence is the same as bisimilarity. In Section 4, we present the generalizations of bisimilarity required to deal with the XPath-algebra and the core XPath-algebra, which are studied in Sections 5 and 6, respectively. The structural characterizations of the semantics of the four XPath fragments in Sections 3, 5 and 6 pertain to the “global view” only. In Section 7, we derive the corresponding characterizations for the “local view.” In Section 8, finally, we discuss some ramifications of our results as well as directions for future research.

Because of space considerations, several proofs are either omitted or only sketched. The proofs of Section 4, many of which require a case analysis, have been moved to an Appendix.

## 2. NOTATION AND TERMINOLOGY

In this paper, *documents* are finite *unordered* node-labeled trees. More formally, a document  $D$  is a 4-tuple  $(V, Ed, r, \lambda)$ , with  $V$  the finite set of nodes,  $Ed \subseteq V \times V$  the set of edges,  $r \in V$  the root and  $\lambda : V \rightarrow \mathcal{L}$  the node-labeling function into an infinite enumerable set  $\mathcal{L}$  of labels.

We next define the fragments of XPath [4] considered in this paper. As observed in the Introduction, we can prune the set of operators considerably, since we are only concerned with (1) *expressibility* on (2) a *single* document.

*Definition 1.* The *XPath-algebra* consists of the primitives  $\varepsilon, \hat{\ell}$  ( $\ell \in \mathcal{L}$ ),  $\emptyset, \downarrow, \uparrow$ , together with the operators.  $E_1/E_2, E_1[E_2], E_1 \cup E_2, E_1 \cap E_2$ , and  $E_1 - E_2$ .

Given a document  $D = (V, Ed, r, \lambda)$ , the *semantics*,  $E(D)$ , of an XPath-algebra expression  $E$  is a binary relation over  $V$ , defined as follows:

- $\varepsilon(D) = \{(n, n) \mid n \in V\}$ ;  $\hat{\ell}(D) = \{(n, n) \mid n \in V \text{ and } \lambda(n) = \ell\}$ ;  $\emptyset(D) = \emptyset$ ;
- $\downarrow(D) = Ed$ ;  $\uparrow(D) = Ed^{-1}$ ;
- $E_1/E_2(D) = \pi_{1,4}\sigma_{2=3}(E_1(D) \times E_2(D))$ ;  $E_1[E_2](D) = \pi_{1,2}\sigma_{2=3}(E_1(D) \times E_2(D))$ ;
- $E_1 \star E_2(D) = E_1(D) \star E_2(D)$ , where “ $\star$ ” stands for “ $\cup$ ”, “ $\cap$ ”, or “ $-$ ”.

Actually, we can show (proof omitted) that the predicate operator “ $E_1[E_2]$ ” is superfluous in the XPath-algebra, but we leave it in because it cannot be omitted in the XPath fragments we define next:

- The *downward XPath-algebra* is the XPath-algebra without “ $\uparrow$ ”.
- The *core XPath-algebra* has the same primitives as the XPath-algebra, together with the operators  $E_1/E_2, E_1[E_2]$  with  $E_2$  a boolean combination<sup>1</sup> of core XPath-algebra expressions, and  $E_1 \cup E_2$ .
- The *downward core XPath-algebra* is the core XPath-algebra without “ $\uparrow$ ”.

Definition 1 reflects the “global” perspective of XPath as working on an entire document, rather than the “local” perspective of XPath as working on a particular node, reflected in Definition 2.

*Definition 2.* Let  $E$  be an XPath-algebra expression, and let  $D = (V, Ed, r, \lambda)$  be a document. For  $m \in V$ ,  $E(D)(m) := \{n \in V \mid (m, n) \in E(D)\}$ .

As the first step in our two-step methodology, we are interested in which nodes in a document we can or cannot distinguish by XPath. Therefore, we define the following equivalence relation:

*Definition 3.* Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . Then  $m_1$  and  $m_2$  are *expression-equivalent* (denoted  $m_1 \equiv_e m_2$ ) if, for each XPath-algebra expression  $E$ ,  $E(D)(m_1) = \emptyset$  if and only if  $E(D)(m_2) = \emptyset$ .

Similarly, we can also define *downward expression equivalence* (denoted as  $m_1 \equiv_{e\downarrow} m_2$ ), *core expression equivalence* (denoted  $m_1 \equiv_{e-} m_2$ ), and *downward core expression equivalence* (denoted  $m_1 \equiv_{e-\downarrow} m_2$ ), each corresponding to one of the XPath-algebra fragments introduced above.

Next, we introduce the notion of *signature* of a pair of a nodes in a document.

*Definition 4.* Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m, n \in V$ . The *signature*  $\text{sig}(m, n)$  is an XPath-algebra expression defined as follows:

1. If  $n$  is a descendant (ancestor) of  $m$ , then  $\text{sig}(m, n) := \downarrow^k$  ( $\text{sig}(m, n) := \uparrow^k$ ), with  $k$  the length of the path between  $m$  and  $n$ .<sup>2</sup>
2. Otherwise, let  $\text{top}(m, n)$  be the least common ancestor of  $m$  and  $n$ . Then

$$\text{sig}(m, n) := \text{sig}(m, \text{top}(m, n)) / \text{sig}(\text{top}(m, n), n).$$

The sequence  $m = p_1, \dots, p_k = n$  of all the intermediate nodes encountered upon computing  $\text{sig}(m, n)(D)(m)$  is called the *path* from  $m$  to  $n$ .

Note that, for  $m_1, m_2, n_1, n_2 \in V$ ,  $(m_2, n_2) \in \text{sig}(m_1, n_1)(D)$  in general does *not* imply that  $\text{sig}(m_1, n_1) = \text{sig}(m_2, n_2)$  unless  $n_1$  is a descendant of  $m_1$ , or vice-versa. For example, in the document  $D$  in Figure 1, *top left*,  $(m_1, m_1) \in \text{sig}(m_1, m_3)$ , while  $\text{sig}(m_1, m_1) = \varepsilon$  and  $\text{sig}(m_1, m_3) = \uparrow^2 / \downarrow^2$ .

We therefore define the following comparison between the signatures of pairs of nodes:

<sup>1</sup>Obtained using union, intersection, and complementation with respect to  $V \times V$ .

<sup>2</sup>The exponent notation denotes repeated composition (“/”). If  $m = n$ , then  $\text{sig}(m, n) := \varepsilon$ .

*Definition 5.* Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2 \in V$ . We say that  $\text{sig}(m_1, n_1) \geq \text{sig}(m_2, n_2)$  if  $(m_2, n_2) \in \text{sig}(m_1, n_1)(D)$ .

We conclude this section with the following observation:

**PROPOSITION 1.** *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2 \in V$ . There exists an XPath-algebra expression  $\text{Sig}(m_1, n_1)$  such that  $(m_2, n_2) \in \text{Sig}(m_1, n_1)(D)$  if and only if  $\text{sig}(m_1, n_1) = \text{sig}(m_2, n_2)$ .*

**PROOF.** If  $n_1$  is a descendant of  $m_1$ , or vice-versa, choosing  $\text{Sig}(m_1, n_1) := \text{sig}(m_1, n_1)$  clearly satisfies all requirements. Otherwise,  $\text{Sig}(m, n) := \text{sig}(m, n) - \text{sig}(\text{parent}(m), \text{parent}(n))$  satisfies all requirements.  $\square$

### 3. CHARACTERIZING THE SEMANTICS OF THE DOWNWARD AND THE DOWNWARD CORE XPATH-ALGEBRAS

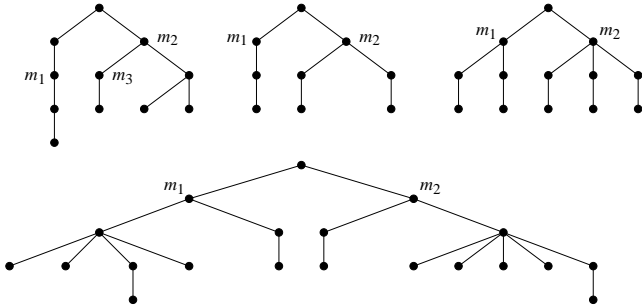
In this section, we are concerned with the downward XPath-algebra and the downward core XPath-algebra, since their semantics have the simplest characterizations. In subsequent sections, we generalize our results to the full XPath-algebra and the core XPath-algebra.

Our first goal is to characterize both downward expression equivalence and downward core expression equivalence in terms of the structure of the document under consideration. Thereto, we define another equivalence relation on the nodes of a document, this time purely in terms of the structure of that document.

*Definition 6.* Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . Then  $m_1$  and  $m_2$  are *downward 1-equivalent* (denoted  $m_1 \equiv_{\downarrow}^1 m_2$ ) if

1.  $\lambda(m_1) = \lambda(m_2)$ ; and
2. for each child  $n_1$  of  $m_1$ , there exists a child  $n_2$  of  $m_2$  such that  $n_1 \equiv_{\downarrow}^1 n_2$ , and vice versa.

In the literature, downward 1-equivalence is usually referred to as *bisimilarity* [5]. For the sake of generalization in Section 4, we use a different terminology in this paper.



**Figure 1: Example documents. All nodes are assumed to have the same label.**

*Example 1.* Consider the document in Figure 1, top left. By Definition 6 the nodes  $m_1$  and  $m_2$  are downward 1-equivalent, whereas the nodes  $m_1$  and  $m_3$  are *not* downward 1-equivalent.

We generalize downward 1-equivalence to *pairs* of nodes.

*Definition 7.* Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2 \in V$  such that  $n_1$  is descendant of  $m_1$  and  $n_2$  is a descendant of  $m_2$ . Then,  $(m_1, n_1)$  and  $(m_2, n_2)$  are *downward 1-equivalent* (denoted  $(m_1, n_1) \equiv_{\downarrow}^1 (m_2, n_2)$ ) if

1.  $\text{sig}(m_1, n_1) = \text{sig}(m_2, n_2)$ ; and
2. for each pair of nodes  $p_1$  and  $p_2$  with
  - (a)  $p_1$  on the path from  $m_1$  to  $n_1$ ;
  - (b)  $p_2$  on the path from  $m_2$  to  $n_2$ ; and
  - (c)  $\text{sig}(m_1, p_1) = \text{sig}(m_2, p_2)$ <sup>3</sup>,

we have that  $p_1 \equiv_{\downarrow}^1 p_2$ .

By repeatedly applying Definition 6, the following connection between downward 1-equivalence of nodes and pairs of nodes can be established:

**LEMMA 1.** *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  such that  $n_1$  is a descendant of  $m_1$  and  $m_1 \equiv_{\downarrow}^1 m_2$ . Then there exists a descendant  $n_2$  of  $m_2$  such that  $(m_1, n_1) \equiv_{\downarrow}^1 (m_2, n_2)$ .*

Using Lemma 1, the following key lemma can now be proved by structural induction.

**LEMMA 2.** *Let  $E$  be a downward XPath-algebra expression, let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2 \in V$  such that  $(m_1, n_1) \equiv_{\downarrow}^1 (m_2, n_2)$ . If  $(m_1, n_1) \in E(D)$ , then  $(m_2, n_2) \in E(D)$ .*

Combining Lemmas 1 and 2 immediately yields

**COROLLARY 1.** *Let  $E$  be a downward XPath-algebra expression, let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  such that  $m_1 \equiv_{\downarrow}^1 m_2$  and  $(m_1, n_1) \in E(D)$ . Then there exists  $n_2 \in V$  such that  $(m_2, n_2) \in E(D)$ .*

We can now present a characterization of downward (core) expression equivalence.

**THEOREM 1.** *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . Then,  $m_1 \equiv_{e\downarrow} m_2$  if and only if  $m_1 \equiv_{e-\downarrow} m_2$  if and only if  $m_1 \equiv_{\downarrow}^1 m_2$ .*

**PROOF.** Since  $m_1 \equiv_{e\downarrow} m_2$  implies  $m_1 \equiv_{e-\downarrow} m_2$ , it remains to prove that (1)  $m_1 \equiv_{\downarrow}^1 m_2$  implies  $m_1 \equiv_{e\downarrow} m_2$  and (2)  $m_1 \equiv_{e-\downarrow} m_2$  implies  $m_1 \equiv_{\downarrow}^1 m_2$ .

For (1), let  $m_1 \equiv_{\downarrow}^1 m_2$ , and let  $E$  be a downward XPath-algebra expression such that  $E(D)(m_1) \neq \emptyset$ . Hence, there exists  $n_1 \in V$  such that  $(m_1, n_1) \in E(D)$ . By Corollary 1, there exists  $n_2 \in V$  such that  $(m_2, n_2) \in E(D)$ , whence  $E(D)(m_2) \neq \emptyset$ . By symmetry, the same holds vice-versa.

For (2), let  $m_1 \equiv_{e-\downarrow} m_2$ . By induction on the height of  $m_1$ , we show that  $m_1 \equiv_{\downarrow}^1 m_2$ .

If  $m_1$  is a leaf, then  $m_2$  is a leaf, for, otherwise,  $\downarrow(D)(m_1) = \emptyset$  and  $\downarrow(D)(m_2) \neq \emptyset$ , a contradiction. In addition, we also have that  $\lambda(m_1) = \lambda(m_2)$ , for, otherwise,  $\widehat{\lambda(m_1)}(D)(m_1) \neq \emptyset$  and  $\widehat{\lambda(m_1)}(D)(m_2) = \emptyset$ , a contradiction. By Definition 6,  $m_1 \equiv_{\downarrow}^1 m_2$ .

If  $m_1$  is not a leaf,  $m_2$  is not a leaf either, and  $\lambda(m_1) = \lambda(m_2)$ , by the same arguments as in the base case. Now, let  $n_1^1$  be a child of  $m_1$ , and let  $n_2^1, \dots, n_2^\ell$  be all children of  $m_2$ . Suppose that,

<sup>3</sup>Or, equivalently,  $\text{sig}(p_1, n_1) = \text{sig}(p_2, n_2)$ .

for all  $i$ ,  $1 \leq i \leq \ell$ ,  $n_1^1 \not\equiv_{e-\downarrow} n_2^i$ . Hence, there exists a downward core XPath-algebra expression  $E_i$  such that  $E_i(D)(n_1^1) \neq \emptyset$  and  $E_i(D)(n_2^i) = \emptyset$ .<sup>4</sup> Let  $F := \varepsilon[\varepsilon[E_1] \cap \dots \cap \varepsilon[E_\ell]]$ . Then  $\downarrow /F(D)(m_1) \neq \emptyset$  and  $\downarrow /F(D)(m_2) = \emptyset$ , a contradiction. Hence, there exists a child  $n_2^j$  of  $m_2$ ,  $1 \leq j \leq \ell$ , such that  $n_1^1 \equiv_{e-\downarrow} n_2^j$ . By the induction hypothesis,  $n_1^1 \equiv_{\downarrow}^1 n_2^j$ . Of course, the same holds vice-versa.  $\square$

As a consequence of Theorem 1, downward (core) expression equivalence is decidable.

We next turn to the second step of our two-step methodology by bootstrapping Theorem 1 to characterize those binary relations over the nodes of a document that can be defined as the evaluation of a downward (core) XPath-algebra expression.<sup>5</sup> For that purpose, we need the following lemma.

**LEMMA 3.** *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . There exists a downward core XPath-algebra expression  $E_{m_1}$  such that  $E_{m_1}(D)(m_2) \neq \emptyset$  if and only if  $m_1 \equiv_{\downarrow}^1 m_2$ .*

**PROOF.** Let  $p_2 \in V$  be a node such that  $m_1 \not\equiv_{\downarrow}^1 p_2$ . By Theorem 1,  $m_1 \not\equiv_{e-\downarrow} p_2$ . Hence, there exists a downward core XPath-algebra expression  $F_{m_1, p_2}$  such that  $F_{m_1, p_2}(D)(m_1) \neq \emptyset$  and  $F_{m_1, p_2}(D)(p_2) = \emptyset$ . It is now easily seen that

$$E_{m_1} := \varepsilon \left[ \bigcap_{p_2 \in V \text{ and } m_1 \not\equiv_{\downarrow}^1 p_2} \varepsilon[F_{m_1 p_2}] \right].$$

is the required downward core XPath-algebra expression.  $\square$

We now prove the main theorem of this section.

**THEOREM 2.** *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $J \subseteq V \times V$ . The following statements are equivalent:*

1. *There exists a core downward XPath-algebra expression  $E$  such that  $E(D) = J$ .*
2. *There exists a downward XPath-algebra expression  $E$  such that  $E(D) = J$ .*
3. (a) *for all  $m, n \in V$ ,  $(m, n) \in J$  implies  $n$  is a descendant of  $m$ ; and*  
 (b) *for all  $m_1, n_1, m_2, n_2 \in V$  with  $n_1$  a descendant of  $m_1$ ,  $n_2$  a descendant of  $m_2$ , and  $(m_1, n_1) \equiv_{\downarrow}^1 (m_2, n_2)$ ,  $(m_1, n_1) \in J$  implies  $(m_2, n_2) \in J$ .*

**PROOF.** Clearly (1)  $\Rightarrow$  (2). The implication (2)  $\Rightarrow$  (3) has been shown in Lemma 2. It remains to show that (3)  $\Rightarrow$  (1). Thereto, consider the downward core XPath-algebra expression

$$E := \bigcup_{(m_1, n_1) \in J} \bigcap_{\substack{p_1 \text{ on the path} \\ \text{from } m_1 \text{ to } n_1}} \text{sig}(m_1, p_1) / \varepsilon[E_{p_1}] / \text{sig}(p_1, n_1),$$

with  $E_{p_1}$  as in Lemma 3. It is now easily seen that condition (3) above implies that  $E(D) = J$ .  $\square$

We immediately conclude that the downward core XPath-algebra and the downward XPath-algebra are equally expressive as navigation tools on a given document.<sup>6</sup>

<sup>4</sup>Alternatively, if  $E'_i$  is an expression such that  $E'_i(D)(n_1^1) = \emptyset$  and  $E'_i(D)(n_2^i) \neq \emptyset$ , then put  $E_i := \varepsilon[\varepsilon - \varepsilon[E'_i]]$ .

<sup>5</sup>In Section 7, we consider this second step for the local view.

<sup>6</sup>Using an involved argument (omitted), we can actually show that both fragments are equivalent as query languages.

## 4. DOWNWARD $k$ -EQUIVALENCE AND $k$ -EQUIVALENCE

We now generalize downward 1-equivalence to downward  $k$ -equivalence, for arbitrary  $k \geq 1$ . The values of  $k$  that will interest us most are 1, 2, and 3.

**Definition 8.** Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . Then  $m_1$  and  $m_2$  are *downward  $k$ -equivalent* (denoted  $m_1 \equiv_{\downarrow}^k m_2$ ) if

1.  $\lambda(m_1) = \lambda(m_2)$ ;
2. for each child  $n_1$  of  $m_1$ , there exists a child  $n_2$  of  $m_2$  such that  $n_1 \equiv_{\downarrow}^k n_2$ , and vice versa; and
3. for each child  $n_1$  of  $m_1$  and each child  $n_2$  of  $m_2$  such that  $n_1 \equiv_{\downarrow}^k n_2$ ,  $\min(|\bar{n}_1|, k) = \min(|\bar{n}_2|, k)$ , where, for  $i = 1, 2$ ,  $\bar{n}_i = \{p \mid (m_i, p) \in Ed \text{ and } p \equiv_{\downarrow}^k n_i\}$ .<sup>7</sup>

Clearly, Definition 8 reduces to Definition 6 for  $k = 1$ . It can be shown (proof omitted) that downward  $k$ -equivalence is the coarsest equivalence relation satisfying conditions (1), (2), and (3) above.

In order to deal with the presence of the “ $\uparrow$ ” operator in both the XPath-algebra and the core XPath-algebra, we need a more restrictive kind of “ $k$ -equivalence” than downward  $k$ -equivalence.

**Definition 9.** Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . Then  $m_1$  and  $m_2$  are  *$k$ -equivalent* (denoted  $m_1 \equiv^k m_2$ ) if

1.  $m_1 \equiv_{\downarrow}^k m_2$ ;
2.  $m_1$  is the root if and only if  $m_2$  is the root;
3. if  $m_1$  and  $m_2$  are not the root, and  $p_1$  and  $p_2$  are the parents of  $m_1$  and  $m_2$ , respectively, then  $p_1 \equiv^k p_2$ .

In other words,  $m_1$  and  $m_2$  are  $k$ -equivalent if they are at the same depth in the document, and each pair of same-generation ancestors of  $m_1$  and  $m_2$  is downward  $k$ -equivalent. As a consequence, we see that same-generation ancestors of  $k$ -equivalent nodes are  $k$ -equivalent themselves.

**Example 2.** In Figure 1, *top left*,  $m_1$  and  $m_2$  are downward 1-equivalent, but *not* 1-equivalent. In Figure 1, *top center*,  $m_1$  and  $m_2$  are 1-equivalent, but *not* 2-equivalent. In Figure 1, *top right*,  $m_1$  and  $m_2$  are 2-equivalent, but *not* 3-equivalent. Finally, in Figure 1, *bottom*,  $m_1$  and  $m_2$  are 3-equivalent, but *not* 4-equivalent.

**Definition 10.** Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2 \in V$ . Then  $(m_1, n_1)$  and  $(m_2, n_2)$  are  *$k$ -equivalent* (denoted  $(m_1, n_1) \equiv^k (m_2, n_2)$ ) if

1.  $\text{sig}(m_1, n_1) = \text{sig}(m_2, n_2)$ ; and
2. for each pair of nodes  $p_1$  and  $p_2$  with
  - (a)  $p_1$  on the path from  $m_1$  to  $n_1$ ;
  - (b)  $p_2$  on the path from  $m_2$  to  $n_2$ ; and
  - (c)  $\text{sig}(m_1, p_1) = \text{sig}(m_2, p_2)$ ,

we have that  $p_1 \equiv^k p_2$ .

Similarly,  $(m_1, n_1)$  and  $(m_2, n_2)$  are  *$k$ -related* (denoted  $(m_1, n_1) \rightleftharpoons^k (m_2, n_2)$ ) if

<sup>7</sup>For a set  $A$ ,  $|A|$  denotes the cardinality of  $A$ .

1.  $\text{sig}(m_1, n_1) \geq \text{sig}(m_2, n_2)$ ; and

2. for each pair of nodes  $p_1$  and  $p_2$  with

- (a)  $p_1$  on the path from  $m_1$  to  $n_1$ ;
- (b)  $p_2$  either on the path from  $m_2$  to  $n_2$  or an ancestor of  $\text{top}(m_2, n_2)$ ; and
- (c)  $\text{sig}(m_1, p_1) \geq \text{sig}(m_2, p_2)$ ,

we have that  $p_1 \equiv^k p_2$ .

Notice that  $k$ -equivalence and  $k$ -relatedness coincide if  $n_1$  is a descendant of  $m_1$ , or vice-versa. In general, downward  $k$ -relatedness is *not* symmetric.

The following technical lemmas are very practical. The second is the generalization of Lemma 1.

**LEMMA 4.** *Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2 \in V$ . Then  $(m_1, n_1) \equiv^k (m_2, n_2)$  (respectively  $(m_1, n_1) \Rightarrow^k (m_2, n_2)$ ) if and only if  $m_1 \equiv^k m_2$ ,  $n_1 \equiv^k n_2$ , and  $\text{sig}(m_1, n_1) = \text{sig}(m_2, n_2)$  (respectively  $\text{sig}(m_1, n_1) \geq \text{sig}(m_2, n_2)$ ).*

**PROOF.** For each pair of nodes  $p_1$  and  $p_2$  for which  $p_1 \equiv^k p_2$  must hold according to Definition 10,  $p_1$  is either an ancestor of  $m_1$  or an ancestor of  $n_1$  and  $p_2$  a same-generation ancestor of  $m_2$  or of  $n_2$ . As observed earlier, same-generation ancestors of  $k$ -equivalent nodes are  $k$ -equivalent.  $\square$

**LEMMA 5.** *Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  such that  $n_1$  is a descendant of  $m_1$  and  $m_1 \equiv^k m_2$ . Then, there exists a descendant  $n_2$  of  $m_2$  such that  $(m_1, n_1) \equiv^k (m_2, n_2)$ .*

The following properties play a crucial role in proving the analogues of Lemma 2 and Corollary 1, used in characterizing the semantics of the downward (core) XPath-algebra, for characterizing the semantics of the XPath-algebra (Lemma 6 and Corollary 2) and the core XPath-algebra (Lemma 11 and Corollary 3). Their proofs are in the Appendix.

**PROPOSITION 2.** *Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  with  $m_1 \equiv^k m_2$ . Then, there exists  $n_2 \in V$  such that  $(m_1, n_1) \Rightarrow^k (m_2, n_2)$ .*

**PROPOSITION 3.** *Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2, p_1 \in V$  such that  $(m_1, n_1) \Rightarrow^k (m_2, n_2)$ . Then, there exists  $p_2 \in V$  such that  $(m_1, p_1) \Rightarrow^k (m_2, p_2)$  and  $(p_1, n_1) \Rightarrow^k (p_2, n_2)$ .*

**PROPOSITION 4.** *Let  $k \geq 2$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  with  $m_1 \equiv^k m_2$ . Then, there exists  $n_2 \in V$  such that  $(m_1, n_1) \equiv^k (m_2, n_2)$ .*

**PROPOSITION 5.** *Let  $k \geq 3$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2, p_1 \in V$  such that  $(m_1, n_1) \equiv^k (m_2, n_2)$ . Then, there exists  $p_2 \in V$  such that  $(m_1, p_1) \equiv^k (m_2, p_2)$  and  $(p_1, n_1) \equiv^k (p_2, n_2)$ .*

## 5. CHARACTERIZING THE SEMANTICS OF THE XPATH-ALGEBRA

Lemma 6, below, is the analogue of Lemma 2 for the full XPath-algebra.

**LEMMA 6.** *Let  $E$  be an XPath-algebra expression, let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2 \in V$  such that  $(m_1, n_1) \equiv^3 (m_2, n_2)$ . If  $(m_1, n_1) \in E(D)$ , then also  $(m_2, n_2) \in E(D)$ .*

**PROOF.** The proof goes by induction on the structure of  $E$ . The induction step for the composition  $E_1/E_2$  relies on Proposition 5; the induction step for the predicate operator  $E_1[E_2]$  relies on Proposition 4; and the induction step for the difference operator  $E_1 - E_2$  relies on the symmetry of 3-equivalence on pairs of nodes. The rest of the proof is straightforward.  $\square$

Combining Proposition 4 and Lemma 6 immediately yields

**COROLLARY 2.** *Let  $E$  be an XPath-algebra expression, let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  such that  $m_1 \equiv^3 m_2$  and  $(m_1, n_1) \in E(D)$ . Then there exists  $n_2 \in V$  such that  $(m_2, n_2) \in E(D)$ .*

Using the same argument used for statement (1) in the proof of Theorem 1, we obtain

**LEMMA 7.** *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . If  $m_1 \equiv^3 m_2$ , then  $m_1 \equiv_e m_2$ .*

The reverse implication, however, requires more work. We first show that expression equivalence implies downward 3-equivalence, and then bootstrap this result to show that, actually, expression equivalence implies 3-equivalence.

**LEMMA 8.** *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . If  $m_1 \equiv_e m_2$ , then  $m_1 \equiv^3 m_2$ .*

**PROOF.** Since downward 3-equivalence is the coarsest equivalence relation satisfying conditions (1), (2), and (3) of Definition 8, it suffices to prove that expression equivalence satisfies these conditions.

For conditions (1) and (2), this requires the same arguments as used for statement (2) in the proof of Theorem 1. We therefore restrict ourselves to condition (3). Thus, let  $n_1^1, \dots, n_k^1$  be all children of  $m_1$  and  $n_2^1, \dots, n_\ell^1$  be all children of  $m_2$ , and assume that  $n_1^1 \equiv_e n_2^1$ . We have to show that  $\min(|\tilde{n}_1^1|, 3) = \min(|\tilde{n}_2^1|, 3)$ , where, for  $i = 1, 2$ ,  $\tilde{n}_i^1 = \{p \mid (m_i^1, p) \in Ed \text{ and } p \equiv_e n_i^1\}$ . To do so, we have to show that the following situations cannot occur:

- 1.  $|\tilde{n}_1^1| = 1$  and  $|\tilde{n}_2^1| > 1$ , or vice-versa; and
- 2.  $|\tilde{n}_1^1| = 2$  and  $|\tilde{n}_2^1| > 2$ , or vice-versa.

By symmetry, it suffices to consider the former situation in each of these cases.

- 1.  $|\tilde{n}_1^1| = 1$  and  $|\tilde{n}_2^1| > 1$ . Hence,  $\tilde{n}_1^1 = \{n_1^1\}$  and, without loss of generality, we may assume that  $\tilde{n}_2^1 \supseteq \{n_2^1, n_2^2\}$ . Since, for all  $i = 2, \dots, k$ ,  $n_1^1 \not\equiv_e n_i^1$ , there exists an XPath-algebra expression  $E_i$  such that  $E_i(D)(n_1^1) \neq \emptyset$  and  $E_i(D)(n_i^1) = \emptyset$ . By definition of expression equivalence, we also have, for  $j = 1, 2$ , that  $E_i(D)(n_j^2) \neq \emptyset$ .

Let  $F := \varepsilon[E_2] \cap \dots \cap \varepsilon[E_k]$ , and let  $G := F / \uparrow / \downarrow / F$ . One can easily verify that  $\varepsilon[G - \varepsilon](D)(n_1^1) = \emptyset$ , while  $\varepsilon[G - \varepsilon](D)(n_2^1) \neq \emptyset$ , a contradiction.<sup>8</sup> So, this case cannot occur.

<sup>8</sup>Of course, one could also have use the expression  $G - \varepsilon$  instead of  $\varepsilon[G - \varepsilon]$ . However, our choice allows reuse of this part of the proof in a subsequent proof.

2.  $|\tilde{n}_1^1| = 2$  and  $|\tilde{n}_2^1| > 2$ . Without loss of generality, we may assume that  $\tilde{n}_1^1 = \{n_1^1, n_2^1\}$  and  $\tilde{n}_2^1 \supseteq \{n_2^1, n_2^2, n_2^3\}$ . Since, for all  $i = 3, \dots, k$ ,  $n_1^1 \neq n_i^1$ , there exists an XPath-algebra expression  $E_i$  such that  $E_i(D)(n_1^1) \neq \emptyset$  and  $E_i(D)(n_i^1) = \emptyset$ . By definition of expression equivalence, we also have, for  $j = 1, 2, 3$ , that  $E_i(D)(n_2^j) \neq \emptyset$ .

Now, let  $F := \varepsilon[E_3] \cap \dots \cap \varepsilon[E_k]$ , let  $G := F / \uparrow / \downarrow / F$ , and let  $H := \varepsilon[G - \varepsilon]$ . One can easily verify that  $((H/H) - \varepsilon)(D)(n_1^1) = \emptyset$ , while  $((H/H) - \varepsilon)(D)(n_2^1) \neq \emptyset$ , a contradiction. So, this case cannot occur either.

We may thus conclude that expression equivalence also satisfies condition (3) of Definition 8.  $\square$

LEMMA 9. *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . If  $m_1 \equiv_e m_2$ , then  $m_1 \equiv^3 m_2$ .*

PROOF. By induction on the depth of  $m_1$  in the document.

If  $m_1$  is the root, then  $m_2$  is also the root, for, otherwise,  $\uparrow(D)(m_1) = \emptyset$  and  $\uparrow(D)(m_2) \neq \emptyset$ . Equal nodes are of course 3-equivalent.

If  $m_1$  is not the root, then  $m_2$  cannot be the root either, for, otherwise, we could derive a contradiction as in the base case. Thus, condition (2) of Definition 9 is met. To prove that condition (3) is met, let  $p_1$  be the parent of  $m_1$  and  $p_2$  the parent of  $m_2$ . If  $p_1 \neq_e p_2$ , there exists an XPath-algebra expression  $E$  such that  $E(D)(p_1) \neq \emptyset$  and  $E(D)(p_2) = \emptyset$ . Obviously, then  $\uparrow/E(D)(m_1) \neq \emptyset$  and  $\uparrow/E(D)(m_2) = \emptyset$ , a contradiction. Thus,  $p_1 \equiv_e p_2$ . By the induction hypothesis,  $p_1 \equiv^3 p_2$ . Finally, Lemma 8 yields condition (1). We may thus conclude that  $m_1 \equiv^3 m_2$ .  $\square$

Lemmas 7 and 9 are both directions of a characterization of expression equivalence:

THEOREM 3. *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . Then,  $m_1 \equiv_e m_2$  if and only if  $m_1 \equiv^3 m_2$ .*

As a consequence of Theorem 3, expression equivalence is decidable.

We next turn to characterizing those binary relations over the nodes of a document that can be defined as the evaluation of an XPath-algebra expression. For that purpose, we need the following lemma, which is the analogue for the full XPath-algebra of Lemma 3 for the downward (core) XPath-algebra. The proof is completely analogous.

LEMMA 10. *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . There exists an XPath-algebra expression  $E_{m_1}$  such that  $E_{m_1}(D)(m_2) \neq \emptyset$  if and only if  $m_1 \equiv^3 m_2$ .*

We now prove the main theorem of this section.

THEOREM 4. *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $J \subseteq V \times V$ . There exists an XPath-algebra expression  $E$  such that  $E(D) = J$  if and only if, for all  $m_1, m_2, n_1, n_2 \in V$  such that  $(m_1, n_1) \equiv^3 (m_2, n_2)$ ,  $(m_1, n_1) \in J$  implies  $(m_2, n_2) \in J$ .*

PROOF. The “only if” follows immediately from Lemma 6. Therefore, we focus on the “if”. Thereto, consider the XPath-algebra expression

$$E := \bigcup_{(m_1, n_1) \in J} \varepsilon[E_{m_1}] / \text{Sig}(m_1, n_1) / \varepsilon[E_{n_1}],$$

with  $E_{m_1}$  and  $E_{n_1}$  as in Lemma 10 and  $\text{Sig}(m_1, n_1)$  as in Proposition 1. It is now easily seen that the condition above imposed on  $J$  implies that  $E(D) = J$ .  $\square$

## 6. CHARACTERIZING THE SEMANTICS OF THE CORE XPath-ALGEBRA

Lemma 11, below, is the analogue of Lemma 6 for the core XPath-algebra.

LEMMA 11. *Let  $E$  be a core XPath-algebra expression, let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2 \in V$  such that  $(m_1, n_1) \equiv^2 (m_2, n_2)$ . If  $(m_1, n_1) \in E(D)$ , then  $(m_2, n_2) \in E(D)$ .*

PROOF. The proof goes by induction on the structure of  $E$ . The proof of the base case is straightforward. The induction step for the composition  $E_1/E_2$  relies on Proposition 3. The induction step for the union operator  $E_1 \cup E_2$  is straightforward. We discuss the induction step for the predicate operator  $E_1[E_2]$ , with  $E_1$  a core XPath-algebra expression and  $E_2$  a boolean combination of core XPath-algebra expressions, in more detail.

Since  $E_2$  can be normalized in disjunctive normal form, and since set union can be pushed out from the predicate to the expression level, we may assume that  $E_2$  is of the form  $F_1 \cap \dots \cap F_k \cap \overline{G_1} \cap \dots \cap \overline{G_\ell}$ . If  $(m_1, n_1) \in E(D)$ , there exists  $p_1 \in V$  such that  $(m_1, n_1) \in E_1(D)$ ,  $(n_1, p_1) \in F_1(D)$ ,  $\dots$ ,  $(n_1, p_1) \in F_k(D)$ ,  $(n_1, p_1) \notin G_1(D)$ ,  $\dots$ ,  $(n_1, p_1) \notin G_\ell(D)$ . By the induction hypothesis,  $(m_2, n_2) \in E_1(D)$ . By Proposition 4, there exists  $p_2 \in V$  such that  $(n_1, p_1) \equiv^2 (n_2, p_2)$ . In particular,  $(n_1, p_1) \equiv^2 (n_2, p_2)$ , whence, by the induction hypothesis,  $(n_2, p_2) \in F_1(D)$ ,  $\dots$ ,  $(n_2, p_2) \in F_k(D)$ . Since  $(n_1, p_1) \equiv^2 (n_2, p_2)$ , we also have  $(n_2, p_2) \equiv^2 (n_1, p_1)$ .<sup>9</sup> If there were  $i$ ,  $1 \leq i \leq \ell$ , such that  $(n_2, p_2) \in G_i(D)$ , then, by the induction hypothesis,  $(n_1, p_1) \in G_i(D)$ , a contradiction. We conclude that  $(n_2, p_2) \notin G_1(D), \dots, (n_2, p_2) \notin G_\ell(D)$ , whence  $(m_2, n_2) \in E(D)$ .  $\square$

Notice that the absence of difference at the expression level is crucial for this proof to work, as an induction step for the difference operator would fail because of the asymmetry of “ $\equiv^2$ ”.

Combining Proposition 2 and Lemma 11 immediately yields

COROLLARY 3. *Let  $E$  be a core XPath-algebra expression, let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  such that  $m_1 \equiv^2 m_2$  and  $(m_1, n_1) \in E(D)$ . Then there exists  $n_2 \in V$  such that  $(m_2, n_2) \in E(D)$ .*

Using the same argument used for statement (1) in the proof of Theorem 1, we obtain

LEMMA 12. *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . If  $m_1 \equiv^2 m_2$ , then  $m_1 \equiv_{e-} m_2$ .*

To prove the reverse direction, we proceed in the same way as for the XPath-algebra.

LEMMA 13. *Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . If  $m_1 \equiv_{e-} m_2$ , then  $m_1 \equiv^2 m_2$ .*

PROOF. The proof is completely analogous to that of Lemma 8, except that, in order to prove that core expression equivalence satisfies condition (3) of Definition 9, we must only show that the case “ $|\tilde{n}_1^1| = 1$  and  $|\tilde{n}_2^1| > 1$ ” cannot occur. Since the expression exhibited for this case is actually a core XPath-algebra expression, the argument used there can be reused here.  $\square$

Notice that the expression exhibited in the proof of Lemma 8 to show that the case “ $|\tilde{n}_1^1| = 2$  and  $|\tilde{n}_2^1| > 2$ ” cannot be transformed into a core XPath-algebra expression.

Lemma 13 can be bootstrapped to Lemma 14, in the same way as Lemma 8 to Lemma 9:

<sup>9</sup>Remember that, while “ $\equiv^2$ ” is symmetric, “ $\equiv^3$ ” in general is not!

LEMMA 14. Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . If  $m_1 \equiv_{e^-} m_2$ , then  $m_1 \equiv^2 m_2$ .

Lemmas 12 and 14 are both directions of a characterization of core expression equivalence:

THEOREM 5. Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . Then,  $m_1 \equiv_{e^-} m_2$  if and only if  $m_1 \equiv^2 m_2$ .

As a consequence of Theorem 5, core expression equivalence is decidable.

We next turn to characterizing those binary relations over the nodes of a document that can be defined as the evaluation of a core XPath-algebra expression. For that purpose, we need the following lemma, which is the analogue for the core XPath-algebra of Lemma 3 for the downward (core) XPath-algebra. The proof is completely analogous.

LEMMA 15. Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2 \in V$ . There exists a core XPath-algebra expression  $E_{m_1}$  such that  $E_{m_1}(D)(m_2) \neq \emptyset$  if and only if  $m_1 \equiv^2 m_2$ .

We now prove the main theorem of this section.

THEOREM 6. Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $J \subseteq V \times V$ . There exists a core XPath-algebra expression  $E$  such that  $E(D) = J$  if and only if, for all  $m_1, m_2, n_1, n_2 \in V$  such that  $(m_1, n_1) \Rightarrow^2 (m_2, n_2)$ ,  $(m_1, n_1) \in J$  implies  $(m_2, n_2) \in J$ .

PROOF. The proof is completely analogous to the proof of Theorem 4, except that, in the expression  $E$  exhibited, “ $\text{Sig}(m_1, n_1)$ ”—which is *not* a core XPath-algebra expression—is replaced by “ $\text{sig}(m_1, n_1)$ ”.  $\square$

## 7. THE LOCAL PERSPECTIVE

Theorems 2, 4, and 6 settle the definability of XPath from a global perspective. Starting from these results, we can now also settle the definability of XPath from a local perspective.

COROLLARY 4. Let  $D = (V, Ed, r, \lambda)$  be a document, let  $m \in V$ , and let  $N \subseteq V$ .

1. There exists a downward (core) XPath-algebra expression  $E$  such that  $E(D)(m) = N$  if and only if, for  $n_1, n_2 \in V$  with  $(m, n_1) \equiv_1^1 (m, n_2)$ ,  $n_1 \in N$  implies  $n_2 \in N$ .
2. There exists an XPath-algebra expression  $E$  where  $E(D)(m) = N$  if and only if, for  $n_1, n_2 \in V$  with  $n_1 \equiv^3 n_2$  and  $\text{sig}(m, n_1) = \text{sig}(m, n_2)$ ,  $n_1 \in N$  implies  $n_2 \in N$ .
3. There exists a core XPath-algebra expression  $E$  such that  $E(D)(m) = N$  if and only if, for  $n_1, n_2 \in V$  with  $n_1 \equiv^2 n_2$  and  $\text{sig}(m, n_1) \geq \text{sig}(m, n_2)$ ,  $n_1 \in N$  implies  $n_2 \in N$ .

For the important special case where the node  $m$  is the root, the statements of Corollary 4 can be simplified.

COROLLARY 5. Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $N \subseteq V$ .

1. There exists a downward (core) XPath-algebra expression  $E$  such that  $E(D)(r) = N$  if and only if, for  $n_1, n_2 \in V$  with  $n_1 \equiv^1 n_2$ ,  $n_1 \in N$  implies  $n_2 \in N$ .
2. There exists an XPath-algebra expression  $E$  where  $E(D)(r) = N$  if and only if, for  $n_1, n_2 \in V$  with  $n_1 \equiv^3 n_2$ ,  $n_1 \in N$  implies  $n_2 \in N$ .
3. There exists a core XPath-algebra expression  $E$  such that  $E(D)(r) = N$  if and only if, for  $n_1, n_2 \in V$  with  $n_1 \equiv^2 n_2$ ,  $n_1 \in N$  implies  $n_2 \in N$ .

## 8. DISCUSSION

In this paper, we characterized the expressive power of four natural fragments of XPath at the document level. Of course, it is possible to consider other fragments or extensions of the XPath-algebra and its data model. Analyzing these using our two-step methodology in order to further improve our understanding of XPath is one possible research direction which we are currently pursuing.

Another future research direction is refining the links between XPath and finite-variable first-order logics [16]. Recently, such links have been established at the level of query semantics. For example, Marx [17, 18] has shown that Core XPath [11] is equivalent to  $\text{FO}_{\text{tree}}^2$ —first-order logic using at most two variables over *ordered* node-labeled trees—interpreted in the signature `child`, `descendant`, and `following_sibling`. Our results establish new links to finite-variable first-order logics at the document level. For example, we can show that, on a given document, the XPath-algebra and  $\text{FO}^3$ —first-order logic with at most three variables—are equivalent in expressive power. Indeed, we can show that, at the document level, the XPath-algebra is equivalent with Tarski’s relation algebra [23] over trees. Tarski and Givant [24] established the link between Tarski’s algebra and  $\text{FO}^3$ . Theorem 3 can then be used to give a new characterization, other than via pebble-games [8, 15], of when two nodes in an unordered tree are indistinguishable in  $\text{FO}^3$ . In this light, connections between other fragments of the XPath-algebra and finite-variable logics must be examined.

The connection between the XPath-algebra and  $\text{FO}^3$  also has ramifications with regard to complexity issues. Indeed, using a result of Grohe [13] which establishes that expression equivalence for  $\text{FO}^3$  is decidable in polynomial time, it follows readily from Theorem 4 and Corollary 4 that the global and local definability problems for the XPath-algebra are decidable in polynomial time. By other arguments, based on the syntactic characterizations in this paper, one can also establish that the global and local definability problems for the other fragments of the XPath-algebra are decidable in polynomial time. As mentioned in the Introduction, this feasibility suggests efficient partitioning and reduction techniques on the set of nodes and the set of paths in a document. Such techniques may be successfully leveraged for various aspects of XML document processing such as indexing, access control, and document compression. This is another research direction which we are currently pursuing.

## 9. ACKNOWLEDGMENTS

We thank Floris Geerts, Jan Hidders, Changqing Lin, Frank Neven, Jan Van den Bussche, and Yuqing Wu for useful discussions. We especially thank Maarten Marx for clarifications about the links between various fragments of XPath and finite-variable logics at the query level.

## 10. REFERENCES

- [1] F. Bancilhon. On the Completeness of Query Languages for Relational Data Bases. In *MFCS*, pages 112–123, Zakopane, Poland, 1978. Springer LNCS 64.
- [2] M. Benedikt, W. Fan, and F. Geerts. XPath Satisfiability in the Presence of DTDs. In *ACM PODS*, pages 25–36, Baltimore, MD, USA, 2005.
- [3] M. Benedikt, W. Fan, and G. M. Kuper. Structural Properties of XPath Fragments. In *ICDT*, pages 79–95, Siena, Italy, 2003. Springer LNCS 2572.
- [4] A. Berglund, S. Boag, D. Chamberlin, M. F. Fernández, M. Kay, J. Robie, and J. Siméon. XML Path Language (XPath) Version 2.0. Technical report, W3C, 2005.



- [5] P. Blackburn, M. de Rijke, and Y. Venema. *Modal Logic*. Cambridge University Press, Cambridge, UK, 2001.
- [6] P. Buneman, M. Grohe, and C. Koch. Path Queries on Compressed XML. In *VLDB*, pages 141–152, Berlin, Germany, 2003.
- [7] A. K. Chandra and D. Harel. Computable Queries for Relational Data Bases. *J. Comp. Sys. Sci.*, 21(2):156–178, 1980.
- [8] H.-D. Ebbinghaus and J. Flum. *Finite Model Theory*. Springer Verlag, Berlin, 1995.
- [9] I. Fundulaki and M. Marx. Specifying Access Control Policies for XML Documents with XPath. In *ACM SACMAT*, pages 61–69, New York, NY, USA, 2004.
- [10] R. Goldman and J. Widom. DataGuides: Enabling Query Formulation and Optimization in Semistructured Databases. In *VLDB*, pages 436–445, Athens, Greece, 1997.
- [11] G. Gottlob and C. Koch. Monadic Queries over Tree-Structured Data. In *IEEE LICS*, pages 189–202, Copenhagen, Denmark, 2002.
- [12] G. Gottlob, C. Koch, and R. Pichler. Efficient Algorithms for Processing XPath Queries. *ACM Trans. Database Syst.*, 30(2):444–491, 2005.
- [13] M. Grohe. Equivalence in Finite-Variable Logics is Complete for Polynomial Time. *Combinatorica*, 19(4):507–532, 1999.
- [14] R. Kaushik, P. Shenoy, P. Bohannon, and E. Gudes. Exploiting Local Similarity for Indexing Paths in Graph-Structured Data. In *IEEE ICDE*, pages 129–140, San Jose, CA, USA, 2002.
- [15] Ł. Krzeczczakowski. Pebble Games on Trees. In *EACSL CSL*, pages 359–371, Vienna, Austria, 2003. Springer LNCS 2803.
- [16] L. Libkin. Logics for Unranked Trees: An Overview. In *EATCS ICALP*, pages 35–50, Lisbon, Portugal, 2005. Springer LNCS 3580.
- [17] M. Marx. Conditional XPath, the First Order Complete XPath Dialect. In *ACM PODS*, pages 13–22, Paris, France, 2004.
- [18] M. Marx and M. de Rijke. Semantic Characterizations of Navigational XPath. *SIGMOD Record*, 34(2):41–46, 2005.
- [19] G. Miklau and D. Suciu. Containment and Equivalence for a Fragment of XPath. *J. ACM*, 51(1):2–45, 2004.
- [20] T. Milo and D. Suciu. Index Structures for Path Expressions. In *ICDT*, pages 277–295, Jerusalem, Israel, 1999.
- [21] J. Paredaens. On the Expressive Power of the Relational Algebra. *Inf. Process. Lett.*, 7(2):107–111, 1978.
- [22] P. Ramanan. Covering Indexes for XML Queries: Bisimulation - Simulation = Negation. In *VLDB*, pages 165–176, Berlin, Germany, 2003.
- [23] A. Tarski. On the Calculus of Relations. *J. Symb. Log.*, 6(3):73–89, 1941.
- [24] A. Tarski and S. Givant. *A Formalization of Set Theory Without Variables*. American Mathematical Society, Providence, RI, USA, 1987.

## APPENDIX

In this Appendix, we provide more details regarding the proofs of Propositions 2–5. For the convenience of the reader, the statements of the results are repeated.

PROPOSITION 2 1. Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  with  $m_1 \equiv^k m_2$ . Then, there exists  $n_2 \in V$  such that  $(m_1, n_1) \equiv^k (m_2, n_2)$ .

PROOF. The case that  $n_1$  is a descendant of  $m_1$  follows from Lemma 5.

If  $n_1$  is not a descendant of  $m_1$ , then let  $t_1 := \text{top}(m_1, n_1)$ . Since  $m_1 \equiv^k m_2$ , Definition 9 implies that there exists an ancestor  $t_2$  of  $m_2$  such that  $t_1 \equiv^k t_2$  and  $\text{sig}(m_1, t_1) = \text{sig}(m_2, t_2)$  (1). By the previous case, there exists a descendant  $n_2$  of  $t_2$  such that  $n_1 \equiv^k n_2$  and  $\text{sig}(t_1, n_1) = \text{sig}(t_2, n_2)$  (2). From (1) and (2), we deduce  $\text{sig}(m_1, n_1) \geq \text{sig}(m_2, n_2)$ . Lemma 4 now yields the desired result.  $\square$

Before proceeding to the proof of Proposition 3, we like to point attention to Figure 2.

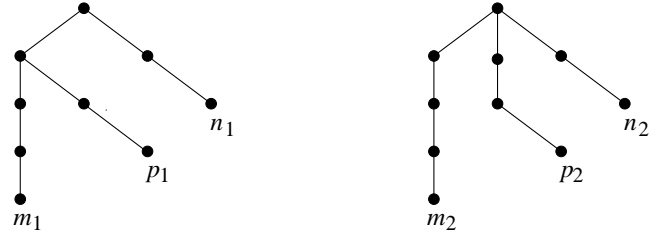


Figure 2: Example document. All nodes are assumed to have the same label

In this document,  $\text{sig}(m_1, n_1) = \text{sig}(m_2, n_2)$  and  $\text{sig}(n_1, p_1) = \text{sig}(n_2, p_2)$ , but  $\text{sig}(m_1, p_1) \not\equiv \text{sig}(m_2, p_2)$ . This example shows that, in the proof of Proposition 3—as well as in the proof of Proposition 5 to follow later—care must be taken in choosing  $p_2$ . Therefore, both proofs proceed via an exhaustive case analysis. Once the correct choice for  $p_2$  is made, however, the remainder of the proof for that case is technical but straightforward and, therefore, omitted.

PROPOSITION 3 1. Let  $k \geq 1$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2, p_1 \in V$  such that  $(m_1, n_1) \equiv^k (m_2, n_2)$ . Then, there exists  $p_2 \in V$  such that  $(m_1, p_1) \equiv^k (m_2, p_2)$  and  $(p_1, n_1) \equiv^k (p_2, n_2)$ .

PROOF. We distinguish three principal cases:

1.  $\text{top}(m_1, p_1)$  is a strict ancestor of  $\text{top}(m_1, n_1)$ .

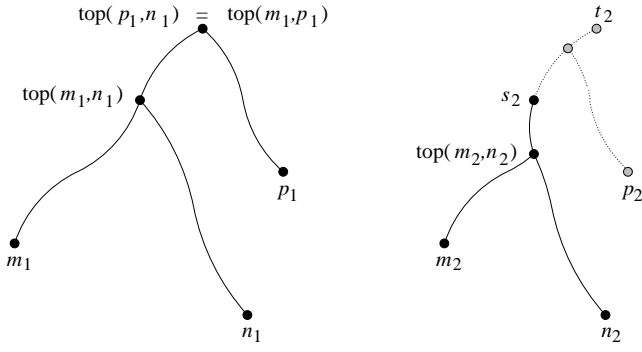
Then,  $\text{top}(p_1, n_1) = \text{top}(m_1, p_1)$ . Let  $p_2$  be any node satisfying  $(m_1, p_1) \equiv^k (m_2, p_2)$ . (Proposition 2). It can now be shown that  $(p_1, n_1) \equiv^k (p_2, n_2)$ . Figure 3 illustrates the constructions in this case for one possible configuration of  $m_2, n_2$ , and  $p_2$ .

2.  $\text{top}(m_1, p_1)$  is a strict descendant of  $\text{top}(m_1, n_1)$ .

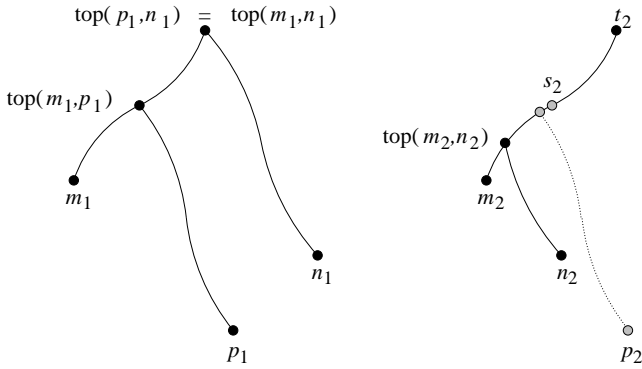
Then,  $\text{top}(p_1, n_1) = \text{top}(m_1, n_1)$ . Now, let  $p_2$  be any node satisfying  $(m_1, p_1) \equiv^k (m_2, p_2)$  (Proposition 2). Again, it can now be shown that  $(p_1, n_1) \equiv^k (p_2, n_2)$ . Figure 4 illustrates the constructions in this case for one possible configuration of  $m_2, n_2$ , and  $p_2$ .

3.  $\text{top}(m_1, p_1) = \text{top}(m_1, n_1)$ . We distinguish two subcases:

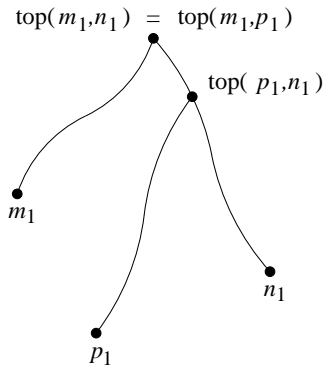
- (a)  $\text{top}(p_1, n_1)$  is a strict descendant of  $\text{top}(m_1, n_1)$ . This situation is shown in Figure 5. This case is the same as the second principal case, with the roles of  $m_1$  and  $n_1$  (and hence the roles of  $m_2$  and  $n_2$ ) interchanged. Since the statement of the lemma is symmetric in this respect, we may consider this case solved.



**Figure 3: Illustration of the constructions in the first principal case of the proof of Proposition 3. The nodes  $s_2$  and  $t_2$  are the ancestors of  $m_2$  satisfying  $\text{top}(m_1, n_1) \equiv^k s_2$  and  $\text{top}(m_1, p_1) \equiv^k t_2$ .**

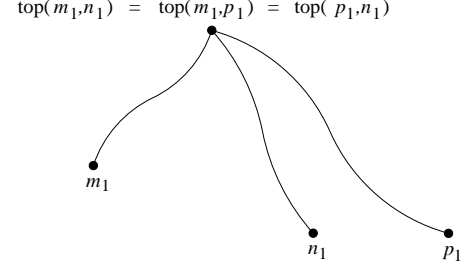


**Figure 4: Illustration of the constructions in the second principal case of the proof of Proposition 3. The nodes  $s_2$  and  $t_2$  are the ancestors of  $m_2$  satisfying  $\text{top}(m_1, n_1) \equiv^k t_2$  and  $\text{top}(m_1, p_1) \equiv^k s_2$ .**



**Figure 5: The first subcase of the third principal case of the proof of Proposition 3.**

(b)  $\text{top}(p_1, n_1) = \text{top}(m_1, p_1) = \text{top}(m_1, n_1)$ . This situation is shown in Figure 6. Since, in this subcase, we have, in particular, that  $\text{top}(m_1, n_1) = \text{top}(p_1, n_1)$  and  $\text{top}(m_1, p_1)$  is on the path from  $m_1$  to  $\text{top}(m_1, n_1)$ , this subcase can be dealt with in the same way as the second principal case.



**Figure 6: The second subcase of the third principal case of the proof of Proposition 3.**

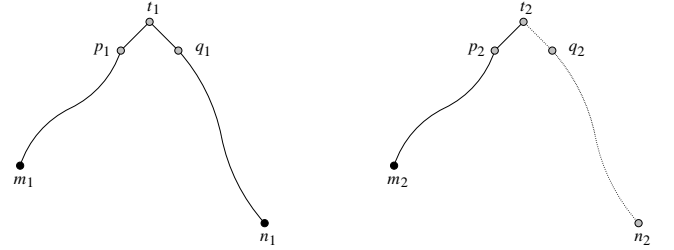
□

**PROPOSITION 4 1.** *Let  $k \geq 2$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1 \in V$  with  $m_1 \equiv^k m_2$ . Then, there exists  $n_2 \in V$  such that  $(m_1, n_1) \equiv^k (m_2, n_2)$ .*

**PROOF.** By Lemma 4, it suffices to show that there exists  $n_2 \in V$  such that  $n_1 \equiv^k n_2$  and  $\text{sig}(m_1, n_1) = \text{sig}(m_2, n_2)$ .

The case where  $n_1$  is a descendant of  $m_1$  is covered by Proposition 2, since  $k$ -equivalence and  $k$ -relatedness coincide in this case.

If  $n_1$  is not a descendant of  $m_1$ , then consider  $t_1 := \text{top}(m_1, n_1)$ . Since  $m_1 \equiv^k m_2$ , Definition 9 implies that there exists an ancestor  $t_2$  of  $m_2$  such that  $t_1 \equiv^k t_2$  and  $\text{sig}(m_1, t_1) = \text{sig}(m_2, t_2)$ .



**Figure 7: Illustration of the constructions in the proof of Proposition 4.**

If  $n_1 = t_1$ , then, clearly,  $n_2 := t_2$  satisfies all requirements. Otherwise, let  $p_1$  be the child of  $t_1$  on the path to  $m_1$ , and let  $q_1$  be the child of  $t_1$  on the path to  $n_1$ . Clearly,  $p_1 \neq q_1$ . Also, let  $p_2$  be the child of  $t_2$  on the path to  $m_2$ . Clearly,  $p_1 \equiv^k p_2$ . In particular,  $p_1 \equiv_{\downarrow}^k p_2$ . We now distinguish two cases:

1.  $p_1 \not\equiv_{\downarrow}^k q_1$ . By Definition 8, there exists a child  $q_2$  of  $t_2$  such that  $q_1 \equiv_{\downarrow}^k q_2$  (whence  $q_1 \equiv^k q_2$ ). Since  $p_1 \not\equiv_{\downarrow}^k q_1$ ,  $p_2 \not\equiv_{\downarrow}^k q_2$ . In particular,  $p_2 \neq q_2$ .
2.  $p_1 \equiv_{\downarrow}^k q_1$ . By Definition 8, and because of  $k \geq 2$ , there exists a child  $q_2$  of  $t_2$  such that  $p_2 \neq q_2$  and  $q_1 \equiv_{\downarrow}^k q_2$  (whence  $q_1 \equiv^k q_2$ ).

In both cases,  $p_2 \neq q_2$  and  $q_1 \equiv^k q_2$ . Since  $n_1$  is a descendant of  $q_1$ , there exists a descendant  $n_2$  of  $q_2$  such that  $n_1 \equiv^k n_2$  and  $\text{sig}(q_1, n_1) = \text{sig}(q_2, n_2)$ . Since  $p_2 \neq q_2$ , it follows that  $\text{sig}(m_1, n_1) = \text{sig}(m_2, n_2)$ . □

PROPOSITION 5 1. Let  $k \geq 3$ . Let  $D = (V, Ed, r, \lambda)$  be a document, and let  $m_1, m_2, n_1, n_2, p_1 \in V$  such that  $(m_1, n_1) \equiv^k (m_2, n_2)$ . Then, there exists  $p_2 \in V$  such that  $(m_1, p_1) \equiv^k (m_2, p_2)$  and  $(p_1, n_1) \equiv^k (p_2, n_2)$ .

PROOF. We distinguish three principal cases:

1.  $top(m_1, p_1)$  is a strict ancestor of  $top(m_1, n_1)$ .

Then,  $top(p_1, n_1) = top(p_2, n_2)$ . Let  $p_2$  be any node satisfying  $(m_1, p_1) \equiv^k (m_2, p_2)$ . (Such a node exists, by Proposition 4.) It is now readily seen that  $(p_1, n_1) \equiv^k (p_2, n_2)$ .

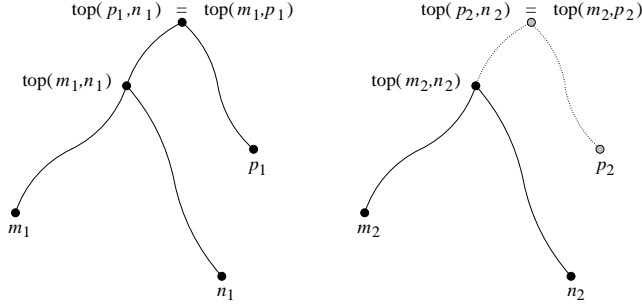


Figure 8: Illustration of the constructions in the first principal case of the proof of Proposition 5.

2.  $top(m_1, p_1)$  is a strict descendant of  $top(m_1, n_1)$ .

Let  $p_2$  be any node satisfying  $(m_1, p_1) \equiv^k (m_2, p_2)$ . (Such a node exists, by Proposition 4.) It is now readily seen that  $(p_1, n_1) \equiv^k (p_2, n_2)$ .

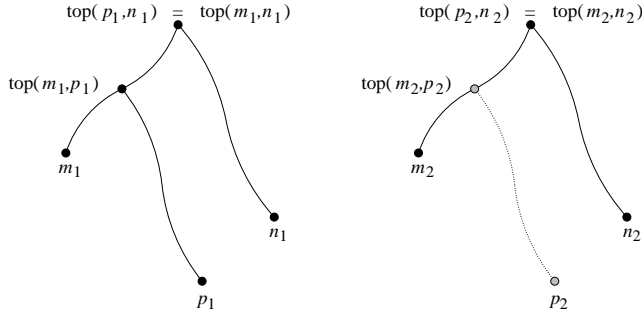


Figure 9: Illustration of the constructions in the second principal case of the proof of Proposition 5.

3.  $top(m_1, p_1) = top(m_1, n_1)$ . We distinguish two subcases:

- (a)  $top(p_1, n_1)$  is a strict descendant of  $top(m_1, n_1)$ .

Let  $p_2$  be any node satisfying  $(p_1, n_1) \equiv^k (p_2, n_2)$ . (Such a node exists, by Proposition 4.) It is now readily seen that  $(m_1, p_1) \equiv^k (m_2, p_2)$ .

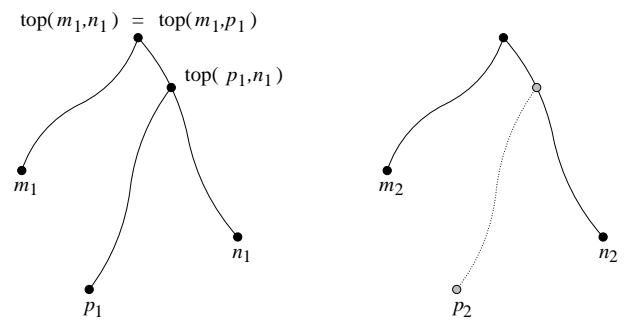


Figure 10: Illustration of the constructions in the first subcase of the third principal case of the proof of Proposition 5.

- (b)  $top(p_1, n_1) = top(m_1, p_1) = top(m_1, n_1)$ .

If  $p_1$  equals this top node, then let  $p_2 := top(m_2, n_2)$ . If  $m_1$  equals this top node, then let  $p_2$  be any node satisfying  $(p_1, n_1) \equiv^k (p_2, n_2)$ . Finally, if  $n_1$  equals this top node, then let  $p_2$  be any node satisfying  $(m_1, p_1) \equiv^k (m_2, p_2)$ . (Such nodes exist, by Proposition 4.) It is readily seen that, in all these border cases,  $p_2$  satisfies all requirements.

If none of these bordercases occur, we are in the situation shown in Figure 11.

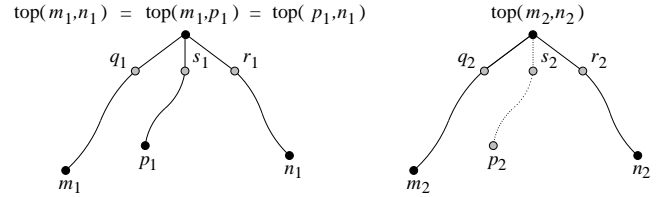


Figure 11: Illustration of the constructions in the second subcase of the third principal case of the proof of Proposition 5.

Let  $q_1, r_1$ , and  $s_1$  be the children of  $top(m_1, n_1)$  on the paths to  $m_1, n_1$ , and  $p_1$ , respectively, and let  $q_2$  and  $r_2$  be the children of  $top(m_2, n_2)$  on the paths to  $m_2$  and  $n_2$ , respectively. Clearly,  $top(m_1, n_1) \equiv^k top(m_2, n_2)$ , whence, in particular,  $top(m_1, n_1) \equiv^k_{\downarrow} top(m_2, n_2)$ . By Definition 8, and since  $k \geq 3$ , it can be seen in an analogous way as in the proof of Proposition 4 that there exists a child  $s_2$  of  $top(m_2, n_2)$  such that  $s_1 \equiv^k_{\downarrow} s_2$  (whence  $s_1 \equiv^k s_2$ ),  $s_2 \neq q_2$ , and  $s_2 \neq r_2$ .

Finally, let  $p_2 \in V$  be any descendant of  $s_2$  satisfying  $(s_1, p_1) \equiv^k (s_2, p_2)$ . (Since  $k$ -equivalence and  $k$ -relatedness coincide for ancestor-descendant pairs, such a node exists, by Proposition 2.) Obviously,  $sig(m_1, p_1) = sig(m_2, p_2)$  and  $sig(p_1, n_1) = sig(p_2, n_2)$ , whence  $(m_1, p_1) \equiv^k (m_2, p_2)$  and  $(p_1, n_1) \equiv^k (p_2, n_2)$ .

□